

## Homework set 5 — MATH 397C — Spring 2022

Due on Thursday May 5. Hand in solutions to problems 1, 2, 6, and 7.

**Problem 1:** The objective of this problem is to computationally investigate the error incurred by truncating multiple expansions. Consider the following geometry: Let  $\Omega_\tau$  and  $\Omega_\sigma$  be two well-separated boxes with centers  $\mathbf{c}_\tau$  and  $\mathbf{c}_\sigma$ . Let  $\mathbf{y} \in \Omega_\sigma$  be a source point and let  $\mathbf{x} \in \Omega_\tau$  be a target point. Consider the error function

$$e(P) = \sup \left\{ \left| \log |\mathbf{x} - \mathbf{y}| - \mathbf{B}_P(\mathbf{x}, \mathbf{c}_\tau) \mathbf{Z}_P(\mathbf{c}_\tau, \mathbf{c}_\sigma) \mathbf{C}_P(\mathbf{c}_\sigma, \mathbf{y}) \right| : \mathbf{y} \in \Omega_\sigma, \mathbf{x} \in \Omega_\tau \right\}$$

where  $P$  is the length of the multipole expansion, and where

$\mathbf{C}_P(\mathbf{c}_\sigma, \mathbf{y}) \in \mathbb{C}^{P \times 1}$	maps a source to an outgoing expansion
$\mathbf{Z}_P(\mathbf{c}_\tau, \mathbf{c}_\sigma) \in \mathbb{C}^{P \times P}$	maps an outgoing expansion to an incoming expansion
$\mathbf{B}_P(\mathbf{x}, \mathbf{c}_\tau) \in \mathbb{C}^{1 \times P}$	maps an incoming expansion to a target

(a) Estimate  $e(P)$  experimentally for the geometry:

$$\Omega_\sigma = [-1, 1] \times [-1, 1], \quad \Omega_\tau = [3, 5] \times [-1, 1].$$

(b) Fit the function you determined in (a) to a curve  $e(P) \sim c \cdot \alpha^P$ . What is  $\alpha$ ?

(c) Is the supremum for a given  $P$  attained for any specific pair  $\{\mathbf{x}, \mathbf{y}\}$ ?

If so, find (experimentally) the pair. Does the choice depend on  $P$ ?

(d) Repeat questions (a), (b), (c) for a different geometry of your choice. (Provide a picture.)

*Hint:* The provided file `main_T_ops_are_fun.m` might be useful.

**Problem 2:** The objective of this exercise is to familiarize yourself with the provided prototype FMM. The questions below refer to the basic FMM provided in the file `main_fmm.m` when executed on a uniform particle distribution. For this case, precompute only the translation operators  $\mathbf{T}^{(\text{of})}$ ,  $\mathbf{T}^{(\text{ifo})}$ , and  $\mathbf{T}^{(\text{if})}$  (i.e. set `flag_precomp=0`).

(a) Estimate and plot the execution time of the FMM for the choices

$$N_{\text{tot}} = 1\,000, 2\,000, 4\,000, 8\,000, 16\,000, 32\,000, 64\,000.$$

Set `nmax=50`. Provide plots that track the following costs:

$t_{\text{tot}}$	total execution time, including initialization.
$t_{\text{init}}$	cost of initialization (computing the tree, the object <code>T_OPS</code> , etc.).
$t_{\text{ofs}}$	cost of applying $\mathbf{T}^{(\text{ofs})}$ .
$t_{\text{of})}$	cost of applying $\mathbf{T}^{(\text{of})}$ .
$t_{\text{ifo}}$	cost of applying $\mathbf{T}^{(\text{ifo})}$ .
$t_{\text{if})}$	cost of applying $\mathbf{T}^{(\text{if})}$ .
$t_{\text{tf})}$	cost of applying $\mathbf{T}^{(\text{tf})}$ .
$t_{\text{close}}$	cost of directly evaluating close range interactions.

(b) Repeat exercise (a) but now for a few different choices of `nmax`. Which one is the best one? Provide a new plot of the times required for this optimal choice.

**Problem 3:** [Optional] Repeat Problem 3.2 but now use a non-uniform point distribution of your choice.

**Problem 4:** [Optional] Can you think of a better way of computing the interaction lists? Here “better” could mean either a cleaner code that executes in more or less the same time, or a code that executes significantly faster than the provided one. If your code is *both* cleaner and faster then so much the better!

**Problem 5:** [Optional] Code up the single-level Barnes-Hut method and investigate computationally how many boxes you should use for optimal performance for any given precision and given total number  $N_{\text{tot}}$  of charges. Create a plot of the best possible time  $t_{\text{optimal}}$  for several  $N_{\text{tot}}$  and estimate the dependence of  $t_{\text{optimal}}$  on  $N_{\text{tot}}$ . To keep things simple, consider only uniform particle distributions. You need only consider a fixed precision (say  $P = 10$ ) but an ambitious solution should compute the optimal time for several different choices (say  $P = 5, 10, 15, 20$ ).

**Problem 6:** In this problem,  $n$  and  $k$  are positive integers such that  $k < n$ ,  $\mathbf{A}$  is an  $N \times N$  invertible matrix, and  $\mathbf{B} = \mathbf{A}^{-1}$ . Let us further assume that every diagonal block of  $\mathbf{A}$  is invertible.

(a) Suppose that  $N = 2n$ , and that we can write  $\mathbf{A}$  and  $\mathbf{B}$  as

$$\mathbf{A} = \begin{bmatrix} \mathbf{A}_{11} & \mathbf{A}_{12} \\ \mathbf{A}_{21} & \mathbf{A}_{22} \end{bmatrix}, \quad \text{and} \quad \mathbf{B} = \begin{bmatrix} \mathbf{B}_{11} & \mathbf{B}_{12} \\ \mathbf{B}_{21} & \mathbf{B}_{22} \end{bmatrix},$$

where each block is of size  $n \times n$ . Suppose further that  $\mathbf{A}_{12}$  and  $\mathbf{A}_{21}$  have rank  $k$ . What is the highest possible value for the rank of  $\mathbf{B}_{12}$ ?

(b) Suppose that  $N = 4n$ , and that we can write  $\mathbf{A}$  and  $\mathbf{B}$  as

$$\mathbf{A} = \begin{bmatrix} \mathbf{A}_{11} & \mathbf{A}_{12} & \mathbf{A}_{13} & \mathbf{A}_{14} \\ \mathbf{A}_{21} & \mathbf{A}_{22} & \mathbf{A}_{23} & \mathbf{A}_{24} \\ \mathbf{A}_{31} & \mathbf{A}_{32} & \mathbf{A}_{33} & \mathbf{A}_{34} \\ \mathbf{A}_{41} & \mathbf{A}_{42} & \mathbf{A}_{43} & \mathbf{A}_{44} \end{bmatrix}, \quad \text{and} \quad \mathbf{B} = \begin{bmatrix} \mathbf{B}_{11} & \mathbf{B}_{12} & \mathbf{B}_{13} & \mathbf{B}_{14} \\ \mathbf{B}_{21} & \mathbf{B}_{22} & \mathbf{B}_{23} & \mathbf{B}_{24} \\ \mathbf{B}_{31} & \mathbf{B}_{32} & \mathbf{B}_{33} & \mathbf{B}_{34} \\ \mathbf{B}_{41} & \mathbf{B}_{42} & \mathbf{B}_{43} & \mathbf{B}_{44} \end{bmatrix},$$

where each block is of size  $n \times n$ . Suppose further that  $\mathbf{A}_{12}, \mathbf{A}_{21}, \mathbf{A}_{34}, \mathbf{A}_{43}, \begin{bmatrix} \mathbf{A}_{13} & \mathbf{A}_{14} \\ \mathbf{A}_{23} & \mathbf{A}_{24} \end{bmatrix}$ , and  $\begin{bmatrix} \mathbf{A}_{31} & \mathbf{A}_{32} \\ \mathbf{A}_{41} & \mathbf{A}_{42} \end{bmatrix}$  all have rank  $k$ . What is the highest possible value for the rank of  $\mathbf{B}_{12}$ ?

(c) [Optional:] Consider the natural generalization to a matrix consisting of  $8 \times 8$  blocks. What is the maximal rank of  $\mathbf{B}_{12}$ ? What about a matrix with  $2^p \times 2^p$  blocks?

Please motivate your answers. If you rely on any formula for the inverse of a blocked matrix, then you may freely assume that any inverse that appears actually exists.

**Problem 7:** In this problem, let  $k$  and  $n$  be integers such that  $0 < k < n$ . Further, let  $\mathbf{D} \in \mathbb{R}^{n \times n}$ ,  $\mathbf{U} \in \mathbb{R}^{n \times k}$ ,  $\mathbf{V} \in \mathbb{R}^{n \times k}$ , and  $\tilde{\mathbf{A}} \in \mathbb{R}^{k \times k}$ .

- (a) Set  $\mathbf{A} = \mathbf{I} + \mathbf{U}\mathbf{V}^*$ . Prove that if  $\mathbf{I} + \mathbf{V}^*\mathbf{U}$  is non-singular, then

$$\mathbf{A}^{-1} = \mathbf{I} - \mathbf{U}(\mathbf{I} + \mathbf{V}^*\mathbf{U})^{-1}\mathbf{V}^*.$$

- (b) Set  $\mathbf{A} = \mathbf{D} + \mathbf{U}\tilde{\mathbf{A}}\mathbf{V}^*$ . Prove that if  $\mathbf{D}$  and  $\mathbf{I} + \tilde{\mathbf{A}}\mathbf{V}^*\mathbf{D}^{-1}\mathbf{U}$  are both invertible, then

$$\mathbf{A}^{-1} = \mathbf{D}^{-1} - \mathbf{D}^{-1}\mathbf{U}(\mathbf{I} + \tilde{\mathbf{A}}\mathbf{V}^*\mathbf{D}^{-1}\mathbf{U})^{-1}\tilde{\mathbf{A}}\mathbf{V}^*\mathbf{D}^{-1}. \quad (1)$$

*Hint:* You could write  $\mathbf{A} = \mathbf{D}(\mathbf{I} + (\mathbf{D}^{-1}\mathbf{U})(\tilde{\mathbf{A}}\mathbf{V}^*))$  and apply (a).

- (c) As in (b), set  $\mathbf{A} = \mathbf{D} + \mathbf{U}\tilde{\mathbf{A}}\mathbf{V}^*$ , and assume that  $\mathbf{D}$  and  $\mathbf{I} + \tilde{\mathbf{A}}\mathbf{V}^*\mathbf{D}^{-1}\mathbf{U}$  are both invertible. Consider the linear system

$$\begin{bmatrix} \mathbf{D} & \mathbf{U}\tilde{\mathbf{A}} \\ -\mathbf{V}^* & \mathbf{I} \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \tilde{\mathbf{x}} \end{bmatrix} = \begin{bmatrix} \mathbf{b} \\ \mathbf{0} \end{bmatrix}. \quad (2)$$

Prove that if  $\begin{bmatrix} \mathbf{x} \\ \tilde{\mathbf{x}} \end{bmatrix}$  solves (2), then  $\mathbf{x}$  solves  $\mathbf{A}\mathbf{x} = \mathbf{b}$ .

- (d) Assume again that  $\mathbf{D}$  is invertible. You can then perform one step of a block LDU factorization to obtain the equality

$$\begin{bmatrix} \mathbf{D} & \mathbf{U}\tilde{\mathbf{A}} \\ -\mathbf{V}^* & \mathbf{I} \end{bmatrix} = \begin{bmatrix} \mathbf{I} & \mathbf{0} \\ \mathbf{X} & \mathbf{I} \end{bmatrix} \begin{bmatrix} \mathbf{D} & \mathbf{0} \\ \mathbf{0} & \mathbf{Y} \end{bmatrix} \begin{bmatrix} \mathbf{I} & \mathbf{Z} \\ \mathbf{0} & \mathbf{I} \end{bmatrix} \quad (3)$$

for some matrices  $\mathbf{X}$ ,  $\mathbf{Y}$ , and  $\mathbf{Z}$ . Specify  $\mathbf{X}$ ,  $\mathbf{Y}$ , and  $\mathbf{Z}$ .

- (e) [Optional:] Observe that the formula (1) is not ideal for inverting a block separable matrix. The reason is that neither  $\tilde{\mathbf{A}}$  nor  $(\mathbf{I} + \tilde{\mathbf{A}}\mathbf{V}^*\mathbf{D}^{-1}\mathbf{U})^{-1}$  is block diagonal, so multiplying them together would be costly. In class, we instead used

$$\mathbf{A}^{-1} = \mathbf{G} + \mathbf{E}(\hat{\mathbf{D}} + \tilde{\mathbf{A}})^{-1}\mathbf{F}^*, \quad (4)$$

where

$$\hat{\mathbf{D}} = (\mathbf{V}^*\mathbf{D}^{-1}\mathbf{U})^{-1},$$

$$\mathbf{E} = \mathbf{D}^{-1}\mathbf{U}\hat{\mathbf{D}},$$

$$\mathbf{F} = (\hat{\mathbf{D}}\mathbf{V}^*\mathbf{D}^{-1})^*,$$

$$\mathbf{G} = \mathbf{D}^{-1} - \mathbf{D}^{-1}\mathbf{U}\hat{\mathbf{D}}\mathbf{V}^*\mathbf{D}^{-1}.$$

For (4) to hold, we need to assume that  $\mathbf{V}^*\mathbf{D}^{-1}\mathbf{U}$  is non-singular. Observe that  $\hat{\mathbf{D}}$ ,  $\mathbf{E}$ ,  $\mathbf{F}$ , and  $\mathbf{G}$  are all block diagonal. Prove (4). *Hint: Huge bonus points if someone can think of a clean and elegant proof. I do not particularly like any that I have thought of, personally.*

*Note: Be on high alert for typos in this problem!!!*