OF MINNESOTA TWIN CITIES

### Numerical Methods in Electronic Structures Calculations

## **Yousef Saad**

Department of Computer Science and Engineering

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#### Current & recent team members team: (CS side)

- Costas Bekas
- Yunkai Zhou
- Susanne Shontz
- Shiv Gowda

- Post-Doc Post-Doc Post-Doc Masters student [Now at NEC, Houston]
- Emmanuel Lorin de La Grandmaison

**Post-doc** [Now in Montreal]

#### What we do:

We develop Numerical Algorithms

- We test them along with materials science researchers
- ▶ We optimize them
- ▶ We install them in materials code (s) e.g., PARSEC
- ▶ .. a few years later ...

#### What we do:

We develop Numerical Algorithms

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▶ We optimize them

▶ We install them in materials code (s) – e.g., PARSEC

▶ .. a few years later ...

We throw them away and repeat

► Algorithms for solving eigenvalue problems

Find alternatives [avoid eigenvectors, eigenvalues]

Solve various related computational problems [TDDFT, computation of dielectric matrix, ...]

### Electronic structures and Schrödinger's equation

Determining matter's electronic structure can be a major challenge:

Number of particules is large [a macroscopic amount contains  $\approx 10^{23}$  electrons and nuclei] and the physical problem is intrinsically complex.

Solution via the many-body Shrödinger equation:  $\frac{H\Psi = E\Psi}{}$ 

▶ In original form the above equation is very complex

#### $\blacktriangleright$ Hamiltonian H is of the form :

$$egin{aligned} egin{aligned} egin{aligned} eta &= -\sum\limits_{i}rac{\hbar^2 mathbf{
aligned}_i^2}{2M_i} - \sum\limits_{j}rac{\hbar^2 mathbf{
aligned}_j^2}{2m} + rac{1}{2}\sum\limits_{i,j}rac{\Sigma_i Z_j e^2}{|ec{R}_i - ec{R}_j|} \ &- \sum\limits_{i,j}rac{Z_i e^2}{|ec{R}_i - ec{r}_j|} + rac{1}{2}\sum\limits_{i,j}rac{e^2}{|ec{r}_i - ec{r}_j|} \end{aligned}$$

 $\Psi = \Psi(r_1, r_2, \dots, r_n, R_1, R_2, \dots, R_N)$  depends on coordinates of all electrons/nuclei.

Involves sums over all electrons / nuclei and their pairs

Note  $\nabla^2 \Psi$  is Laplacean of  $\Psi$  = sum of second derivatives of  $\Psi$  in each direction. Represents kinetic energy.

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A hypothetical calculation: (dont try this at home) 10 Atoms each having 14 electrons [Silicon] ▶ ... a total of 15\*10= 150 particles Assume each coordinate will need 100 points for discretization.. ▶ ... you will get **# Unknowns** =  $\underline{100} \times \underline{100} \times \cdots \times \underline{100} = 100^{150}$ part.1 part.2 part.150Methods based on this basic formulation are limited to

a few atoms.

### Several approximations/ theories used

Problem can be viewed from the angle of optimization: Find  $\Psi$  Minimize energy  $\leq \Psi \mid H \mid \Psi >$  $\leq \Psi \mid \Psi \mid \Psi >$ 

.. or from the angle of eigenvalue problems: Find eigenfunction  $\Psi$  associated with smallest eigenvalue of H

Methods have been developed on both camps

► Our camp: eigenvalue problems

### A brief history of diagonalization methods

First: what is an eigenvalue problem? Given a matrix A, find a scalar  $\lambda$  and a nonzero vector x such that  $Ax = \lambda x$ 

**Quiz:** what are the main uses of eigenvalues/ eigenvectors

### **Eigenvalue Problems. Their origins**

- Structural Engineering [ $Ku = \lambda Mu$ ]
- Stability analysis [e.g., electrical networks, mechanical system,..]
- Bifurcation analysis [e.g., in fluid flow]
- Electronic structure calculations [Shrödinger equation..]
- Application of new era: page ranking on the world-wide web.

#### **Types of Problems:**

- Compute a few  $\lambda_i$  's with smallest or largest real parts;
- Compute all  $\lambda_i$ 's in a certain region of  $\mathbb{C}$ ;
- Compute a few of the dominant eigenvalues;
- Compute all  $\lambda_i$ 's.

**Our problem:** A is symmetric real – so its eigenvalues are real, and we want to compute the lowest k, where k represents the number of occupied states.

 $\blacktriangleright$  k can be in the hundreds or thousands

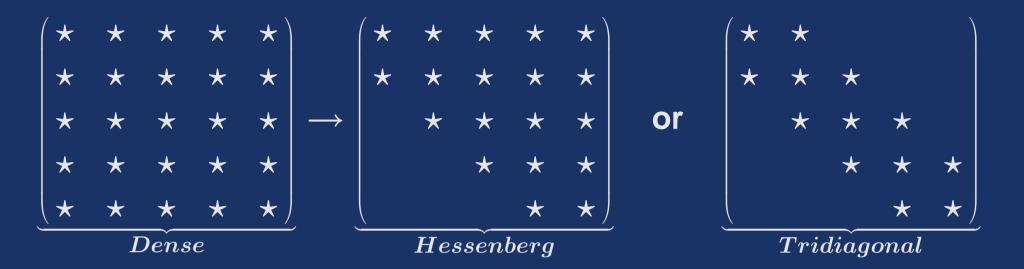
 $\blacktriangleright$  A can be very large up to a few Millions - currently.

#### **Eigenvalues of dense matrices**

**Early days: compute the characteristic polynomial.** 

- Works well "by hand" for dimension of up to 5 or so..
- ► Not viable for serious calculations

▶ Work in the 40's and 50s concentrated on reducing the problem into an easy one: triadiagonal / or Hessenberg form (nonsymmetric matrices).



#### Characteristic polynomials easier to compute...

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### **Breakthroughs: The LR and QR algorithms**

▶ Discovered in 1959 (LR) and 1961 (QR).

- Complicated to understand theoretically
- QR was held "secret" for a few years given its importance.
- Both QR and LR are very economical for tridiagonal matrices.

#### A brief history of eigenvalues - cont.

▶ 1965: Major book in computing eigenvalues by J. Wilkison: The "bible" on eigenvalue problems

▶ 1971: volume by Wilkinson and Reinsch "handbook for automatic computations", Linear algebra. Published 1st programs – in Algol... a now defunct language

► By 1975: EISPACK project. Translate Wilkinson & Reinsch volume in FORTRAN IV. Tremendous impact

Later came LINPACK [linear equations]

▶ and much later came LAPACK = combined the two..

**Current state of the art for dense computations:** 

- **Consider the symmetric case only. Then**
- $\blacktriangleright$  Reduce A to tridiagonal form
- ► Use QR-alg. on tridiag matrix.
- $\blacktriangleright$  Order  $n^3$  calculation –
- **software: LAPACK**

Quiz: Suppose cost is  $14n^3$  and  $n = 10^6$ , and you have a very powerful machine at home which can store the matrix and which delivers an operation every nanosecond (1 Gflop machine). How long would it take to compute the eigenvalues of A? Quiz: Suppose cost is  $14n^3$  and  $n = 10^6$ , and you have a very powerful machine at home which can store the matrix and which delivers an operation every nanosecond (1 Gflop machine). How long would it take to compute the eigenvalues of A?





### Effective methods exploit sparsity

Most methods exploit the fact that a product of a sparse matrix by a vector is very inexpensive [Order *n*]

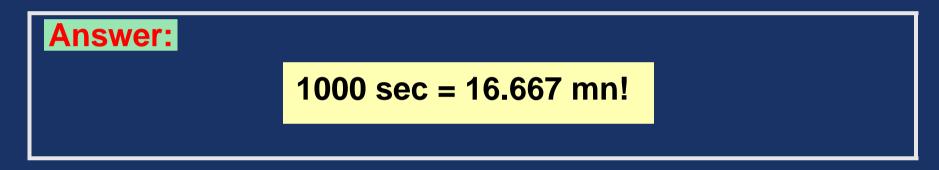
Suppose for the sake of argument that we now have an algorithm that takes exactly  $n^2$  operations to compute one eigenvector.

**Same question as before for one eigenvector** 

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### Environments used throughout the project:

- Initially: Cray YMP [93]
- Cluster of SGI workstations
- CM5 [1994-1996]
- IBM SP2 [Using PVM]
- Cray T3D [Combining PVM + MPI] around 1996-1997
- Cray T3E [using MPI] 1997
- IBM SP with +256 nodes 1998+
- IBM SP3 / SGI Altix currently