

**Input:** An  $m \times n$  matrix  $\mathbf{A}$  and a target rank  $k$ .

**Output:** Rank- $k$  factors  $\mathbf{U}$ ,  $\mathbf{D}$ , and  $\mathbf{V}$  in an approximate SVD  $\mathbf{A} \approx \mathbf{UDV}^*$ .

(1) Draw an  $n \times k$  **random matrix**  $\mathbf{R}$ .

(2) Form the  $m \times k$  **sample matrix**  $\mathbf{Y} = \mathbf{AR}$ .

(3) Compute an **ON matrix**  $\mathbf{Q}$  s.t.  $\mathbf{Y} = \mathbf{QQ}^*\mathbf{Y}$ .

(4) Form the small matrix  $\mathbf{B} = \mathbf{Q}^* \mathbf{A}$ .

(5) Factor the small matrix  $\mathbf{B} = \hat{\mathbf{U}}\mathbf{D}\mathbf{V}^*$ .

(6) Form  $\mathbf{U} = \mathbf{Q}\hat{\mathbf{U}}$ .

**Question:** What is the error  $e_k = \|\mathbf{A} - \mathbf{UDV}^*\|$ ? (Recall that  $e_k = \|\mathbf{A} - \mathbf{QQ}^*\mathbf{A}\|$ .)

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**Answer:** Lamentably, no. The expectation of  $\frac{e_k}{\sigma_{k+1}}$  is large, and has very large variance.

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**Question:** Is  $e_k$  close to  $\sigma_{k+1}$ ?

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**Remedy:** Over-sample *slightly*. Compute  $k+p$  samples from the range of  $\mathbf{A}$ .

It turns out that  $p = 5$  or  $10$  is often sufficient.  $p = k$  is almost always more than enough.

**Input:** An  $m \times n$  matrix  $\mathbf{A}$ , a target rank  $k$ , **and an over-sampling parameter  $p$  (say  $p = 5$ )**.

**Output:** Rank- $(k + p)$  factors  $\mathbf{U}$ ,  $\mathbf{D}$ , and  $\mathbf{V}$  in an approximate SVD  $\mathbf{A} \approx \mathbf{UDV}^*$ .

(1) Draw an  $n \times (k + p)$  **random matrix**  $\mathbf{R}$ .

(2) Form the  $m \times (k + p)$  **sample matrix**  $\mathbf{Y} = \mathbf{AR}$ .

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## ***Bound on the expectation of the error for Gaussian test matrices***

Let  $\mathbf{A}$  denote an  $m \times n$  matrix with singular values  $\{\sigma_j\}_{j=1}^{\min(m,n)}$ .

Let  $k$  denote a target rank and let  $p$  denote an over-sampling parameter.

Let  $\mathbf{R}$  denote an  $n \times (k + p)$  Gaussian matrix.

Let  $\mathbf{Q}$  denote the  $m \times (k + p)$  matrix  $\mathbf{Q} = \text{orth}(\mathbf{AR})$ .

If  $p \geq 2$ , then

$$\mathbb{E} \|\mathbf{A} - \mathbf{Q}\mathbf{Q}^* \mathbf{A}\|_{\text{Frob}} \leq \left(1 + \frac{k}{p-1}\right)^{1/2} \left(\sum_{j=k+1}^{\min(m,n)} \sigma_j^2\right)^{1/2},$$

and

$$\mathbb{E} \|\mathbf{A} - \mathbf{Q}\mathbf{Q}^* \mathbf{A}\| \leq \left(1 + \sqrt{\frac{k}{p-1}}\right) \sigma_{k+1} + \frac{e \sqrt{k+p}}{p} \left(\sum_{j=k+1}^{\min(m,n)} \sigma_j^2\right)^{1/2}.$$

*Ref: Halko, Martinsson, Tropp, 2009 & 2011*

## Large deviation bound for the error for Gaussian test matrices

Let  $\mathbf{A}$  denote an  $m \times n$  matrix with singular values  $\{\sigma_j\}_{j=1}^{\min(m,n)}$ .

Let  $k$  denote a target rank and let  $p$  denote an over-sampling parameter.

Let  $\mathbf{R}$  denote an  $n \times (k + p)$  Gaussian matrix.

Let  $\mathbf{Q}$  denote the  $m \times (k + p)$  matrix  $\mathbf{Q} = \text{orth}(\mathbf{AR})$ .

If  $p \geq 4$ , and  $u$  and  $t$  are such that  $u \geq 1$  and  $t \geq 1$ , then

$$\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| \leq \left(1 + t \sqrt{\frac{3k}{p+1}} + ut \frac{e \sqrt{k+p}}{p+1}\right) \sigma_{k+1} + \frac{te \sqrt{k+p}}{p+1} \left(\sum_{j>k} \sigma_j^2\right)^{1/2}$$

except with probability at most  $2t^{-p} + e^{-u^2/2}$ .

*Ref: Halko, Martinsson, Tropp, 2009 & 2011; Martinsson, Rokhlin, Tygert (2006)*

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$u$  and  $t$  parameterize “bad” events — large  $u$ ,  $t$  is bad, but unlikely.

Certain choices of  $t$  and  $u$  lead to simpler results. For instance,

$$\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| \leq \left(1 + 16\sqrt{1 + \frac{k}{p+1}}\right) \sigma_{k+1} + 8 \frac{\sqrt{k+p}}{p+1} \left(\sum_{j>k} \sigma_j^2\right)^{1/2},$$

except with probability at most  $3e^{-p}$ .

## Large deviation bound for the error for Gaussian test matrices

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If  $p \geq 4$ , and  $u$  and  $t$  are such that  $u \geq 1$  and  $t \geq 1$ , then

$$\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| \leq \left(1 + t \sqrt{\frac{3k}{p+1}} + ut \frac{e \sqrt{k+p}}{p+1}\right) \sigma_{k+1} + \frac{te \sqrt{k+p}}{p+1} \left(\sum_{j>k} \sigma_j^2\right)^{1/2}$$

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$u$  and  $t$  parameterize “bad” events — large  $u$ ,  $t$  is bad, but unlikely.

Certain choices of  $t$  and  $u$  lead to simpler results. For instance,

$$\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| \leq \left(1 + 6\sqrt{(k+p) \cdot p \log p}\right) \sigma_{k+1} + 3\sqrt{k+p} \left(\sum_{j>k} \sigma_j^2\right)^{1/2},$$

except with probability at most  $3p^{-p}$ .



## Proofs — Overview:

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Let  $\mathbf{A}$  denote an  $m \times n$  matrix with singular values  $\{\sigma_j\}_{j=1}^{\min(m,n)}$ .

Let  $k$  denote a target rank and let  $p$  denote an over-sampling parameter. Set  $\ell = k + p$ .

Let  $\mathbf{R}$  denote an  $n \times \ell$  “test matrix”, and let  $\mathbf{Q}$  denote the  $m \times \ell$  matrix  $\mathbf{Q} = \text{orth}(\mathbf{A}\mathbf{R})$ .

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We seek to bound the error  $e_k = e_k(\mathbf{A}, \mathbf{R}) = \|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|$ , which is a random variable.

1. Make no assumption on  $\mathbf{R}$ . Construct a deterministic bound of the form

$$\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| \leq \dots \mathbf{A} \dots \mathbf{R} \dots$$

2. Assume that  $\mathbf{R}$  is drawn from a normal Gaussian distribution.

Take expectations of the deterministic bound to attain a bound of the form

$$\mathbb{E}[\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|] \leq \dots \mathbf{A} \dots$$

3. Assume that  $\mathbf{R}$  is drawn from a normal Gaussian distribution.

Take expectations of the deterministic bound conditioned on “bad behavior” in  $\mathbf{R}$  to get that

$$\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| \leq \dots \mathbf{A} \dots$$

holds with probability at least  $\dots$ .

## Part 1 (out of 3) — deterministic bound:

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Let  $\mathbf{A}$  denote an  $m \times n$  matrix with singular values  $\{\sigma_j\}_{j=1}^{\min(m,n)}$ .

Let  $k$  denote a target rank and let  $p$  denote an over-sampling parameter. Set  $\ell = k + p$ .

Let  $\mathbf{R}$  denote an  $n \times \ell$  “test matrix”, and let  $\mathbf{Q}$  denote the  $m \times \ell$  matrix  $\mathbf{Q} = \text{orth}(\mathbf{A}\mathbf{R})$ .

---

Partition the SVD of  $\mathbf{A}$  as follows:

$$\mathbf{A} = \mathbf{U} \begin{bmatrix} \mathbf{D}_1 & \\ & \mathbf{D}_2 \end{bmatrix} \begin{bmatrix} \mathbf{V}_1^* \\ \mathbf{V}_2^* \end{bmatrix} \begin{matrix} k \\ n - k \end{matrix}$$

Define  $\mathbf{R}_1$  and  $\mathbf{R}_2$  via

$$\begin{matrix} \mathbf{R}_1 & = & \mathbf{V}_1^* & \mathbf{R} \\ k \times (k + p) & & k \times n & n \times (k + p) \end{matrix} \quad \text{and} \quad \begin{matrix} \mathbf{R}_2 & = & \mathbf{V}_2^* & \mathbf{R} \\ (n - k) \times (k + p) & & (n - k) \times n & n \times (k + p) \end{matrix}$$

**Theorem:** [HMT2009,HMT2011] Assuming that  $\mathbf{R}_1$  is not singular, it holds that

$$\|\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|\|^2 \leq \underbrace{\|\|\mathbf{D}_2\|\|^2}_{\text{theoretically minimal error}} + \|\|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\|\|^2.$$

Here,  $\|\|\cdot\|\|$  represents either  $\ell^2$ -operator norm, or the Frobenius norm.

*Note:* A similar (but weaker) result appears in Boutsidis, Mahoney, Drineas (2009).

**Recall:**  $\mathbf{A} = \mathbf{U} \begin{bmatrix} \mathbf{D}_1 & \mathbf{0} \\ \mathbf{0} & \mathbf{D}_2 \end{bmatrix} \begin{bmatrix} \mathbf{V}_1^* \\ \mathbf{V}_2^* \end{bmatrix}$ ,  $\begin{bmatrix} \mathbf{R}_1 \\ \mathbf{R}_2 \end{bmatrix} = \begin{bmatrix} \mathbf{V}_1^* \mathbf{R} \\ \mathbf{V}_2^* \mathbf{R} \end{bmatrix}$ ,  $\mathbf{Y} = \mathbf{A}\mathbf{R}$ ,  $\mathbf{P}$  proj<sup>n</sup> onto  $\text{Ran}(\mathbf{Y})$ .

**Thm:** Suppose  $\mathbf{D}_1 \mathbf{R}_1$  has full rank. Then  $\|\mathbf{A} - \mathbf{P}\mathbf{A}\|^2 \leq \|\mathbf{D}_2\|^2 + \|\mathbf{D}_2 \mathbf{R}_2 \mathbf{R}_1^\dagger\|^2$ .

**Proof:** The problem is rotationally invariant  $\Rightarrow$  We can assume  $\mathbf{U} = \mathbf{I}$  and so  $\mathbf{A} = \mathbf{D}\mathbf{V}^*$ .

Simple calculation:  $\|(\mathbf{I} - \mathbf{P})\mathbf{A}\|^2 = \|\mathbf{A}^*(\mathbf{I} - \mathbf{P})^2\mathbf{A}\| = \|\mathbf{D}(\mathbf{I} - \mathbf{P})\mathbf{D}\|$ .

$$\text{Ran}(\mathbf{Y}) = \text{Ran} \left( \begin{bmatrix} \mathbf{D}_1 \mathbf{R}_1 \\ \mathbf{D}_2 \mathbf{R}_2 \end{bmatrix} \right) = \text{Ran} \left( \begin{bmatrix} \mathbf{I} \\ \mathbf{D}_2 \mathbf{R}_2 \mathbf{R}_1^\dagger \mathbf{D}_1 \end{bmatrix} \mathbf{D}_1 \mathbf{R}_1 \right) = \text{Ran} \left( \begin{bmatrix} \mathbf{I} \\ \mathbf{D}_2 \mathbf{R}_2 \mathbf{R}_1^\dagger \mathbf{D}_1 \end{bmatrix} \right)$$

Set  $\mathbf{F} = \mathbf{D}_2 \mathbf{R}_2 \mathbf{R}_1^\dagger \mathbf{D}_1^{-1}$ . Then  $\mathbf{P} = \begin{bmatrix} \mathbf{I} \\ \mathbf{F} \end{bmatrix} (\mathbf{I} + \mathbf{F}^* \mathbf{F})^{-1} [\mathbf{I} \ \mathbf{F}^*]$ . (Compare to  $\mathbf{P}_{\text{ideal}} = \begin{bmatrix} \mathbf{I} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} \end{bmatrix}$ .)

Use properties of psd matrices:  $\mathbf{I} - \mathbf{P} \preceq \dots \preceq \begin{bmatrix} \mathbf{F}^* \mathbf{F} & -(\mathbf{I} + \mathbf{F}^* \mathbf{F})^{-1} \mathbf{F}^* \\ -\mathbf{F}(\mathbf{I} + \mathbf{F}^* \mathbf{F})^{-1} & \mathbf{I} \end{bmatrix}$

Conjugate by  $\mathbf{D}$  to get  $\mathbf{D}(\mathbf{I} - \mathbf{P})\mathbf{D} \preceq \begin{bmatrix} \mathbf{D}_1 \mathbf{F}^* \mathbf{F} \mathbf{D}_1 & -\mathbf{D}_1 (\mathbf{I} + \mathbf{F}^* \mathbf{F})^{-1} \mathbf{F}^* \mathbf{D}_2 \\ -\mathbf{D}_2 \mathbf{F} (\mathbf{I} + \mathbf{F}^* \mathbf{F})^{-1} \mathbf{D}_1 & \mathbf{D}_2^2 \end{bmatrix}$

Diagonal dominance:  $\|\mathbf{D}(\mathbf{I} - \mathbf{P})\mathbf{D}\| \leq \|\mathbf{D}_1 \mathbf{F}^* \mathbf{F} \mathbf{D}_1\| + \|\mathbf{D}_2^2\| = \|\mathbf{D}_2 \mathbf{R}_2 \mathbf{R}_1^\dagger\|^2 + \|\mathbf{D}_2\|^2$ .

## Part 2 (out of 3) — bound on expectation of error when $\mathbf{R}$ is Gaussian:

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Let  $\mathbf{A}$  denote an  $m \times n$  matrix with singular values  $\{\sigma_j\}_{j=1}^{\min(m,n)}$ .

Let  $k$  denote a target rank and let  $p$  denote an over-sampling parameter. Set  $\ell = k + p$ .

Let  $\mathbf{R}$  denote an  $n \times \ell$  “test matrix”, and let  $\mathbf{Q}$  denote the  $m \times \ell$  matrix  $\mathbf{Q} = \text{orth}(\mathbf{A}\mathbf{R})$ .

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**Recall:**  $\|\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|\|^2 \leq \|\|\mathbf{D}_2\|\|^2 + \|\|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\|\|^2$ , where  $\mathbf{R}_1 = \mathbf{V}_1^*\mathbf{R}$  and  $\mathbf{R}_2 = \mathbf{V}_2^*\mathbf{R}$ .

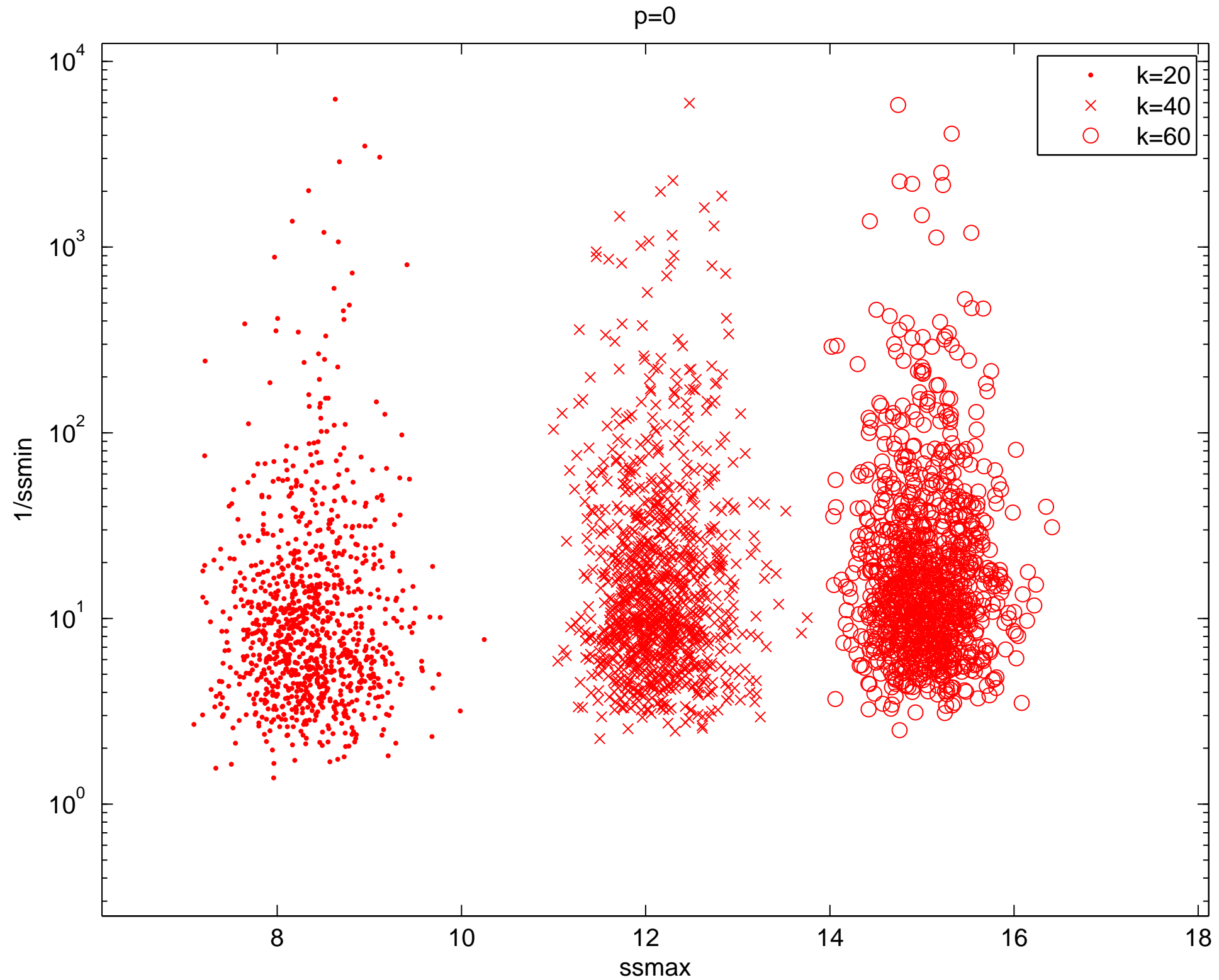
**Assumption:**  $\mathbf{R}$  is drawn from a normal Gaussian distribution.

Since the Gaussian distribution is rotationally invariant, the matrices  $\mathbf{R}_1$  and  $\mathbf{R}_2$  also have a Gaussian distribution. (As a consequence, the matrices  $\mathbf{U}$  and  $\mathbf{V}$  do not enter the analysis and one could simply assume that  $\mathbf{A}$  is diagonal,  $\mathbf{A} = \text{diag}(\sigma_1, \sigma_2, \dots)$ .)

What is the distribution of  $\mathbf{R}_1^\dagger$  when  $\mathbf{R}_1$  is a  $k \times (k + p)$  Gaussian matrix?

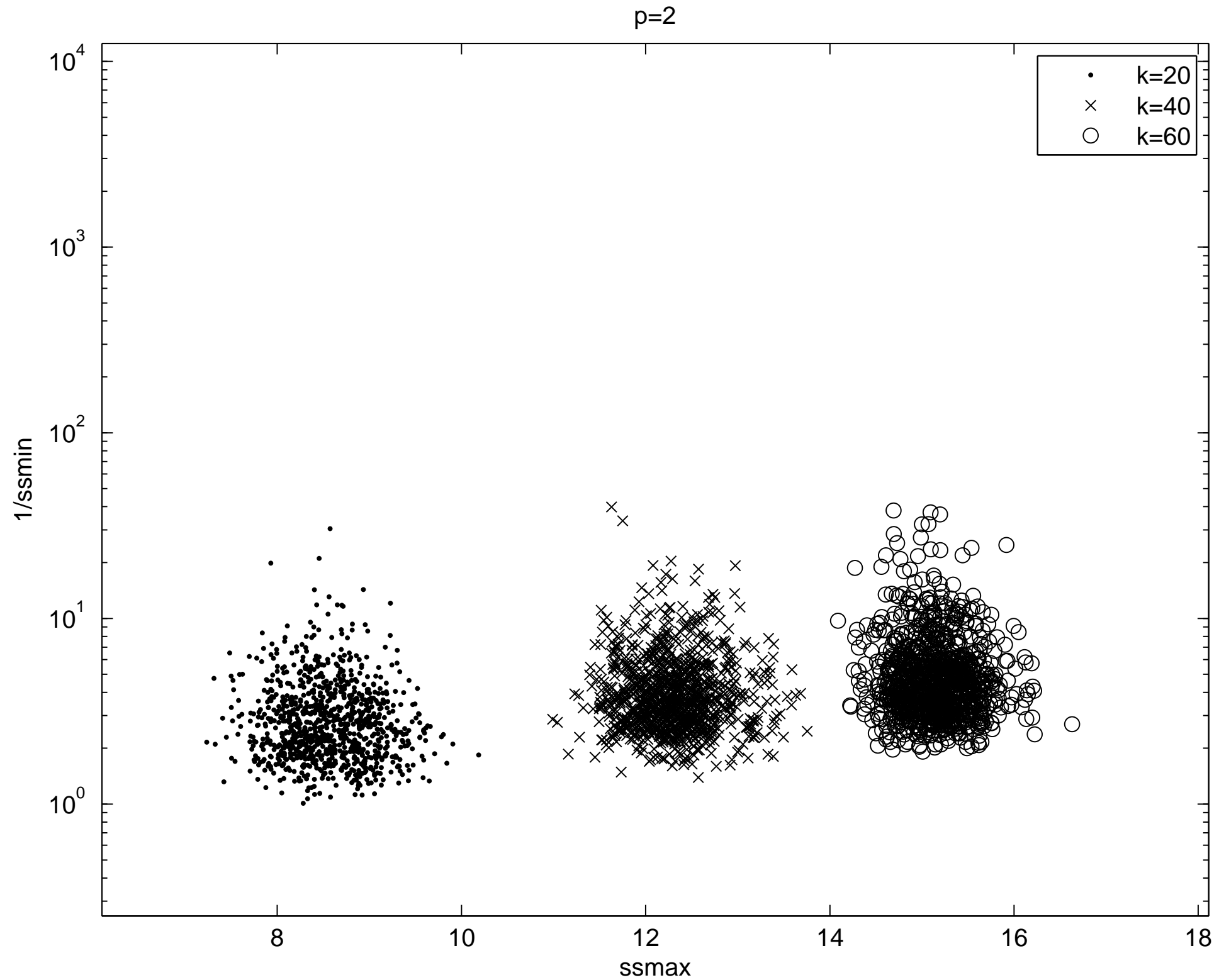
If  $p = 0$ , then  $\|\|\mathbf{R}_1^\dagger\|\|$  is typically large, and is very unstable.

Scatter plot showing distribution of  $1/\sigma_{\min}$  for  $k \times (k + p)$  Gaussian matrices.  $p = 0$



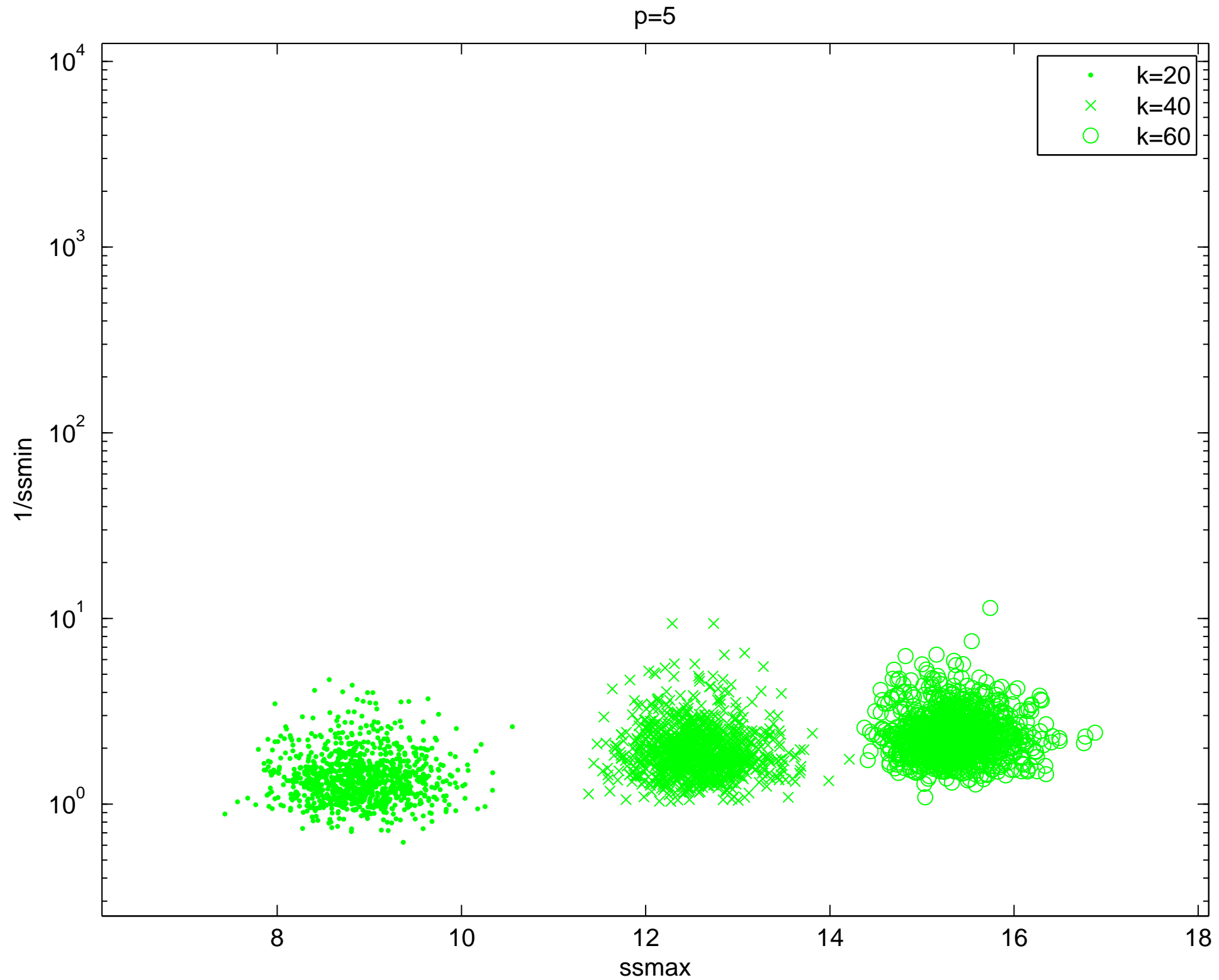
$1/\sigma_{\min}$  is plotted against  $\sigma_{\max}$ .

Scatter plot showing distribution of  $1/\sigma_{\min}$  for  $k \times (k + p)$  Gaussian matrices.  $p = 2$



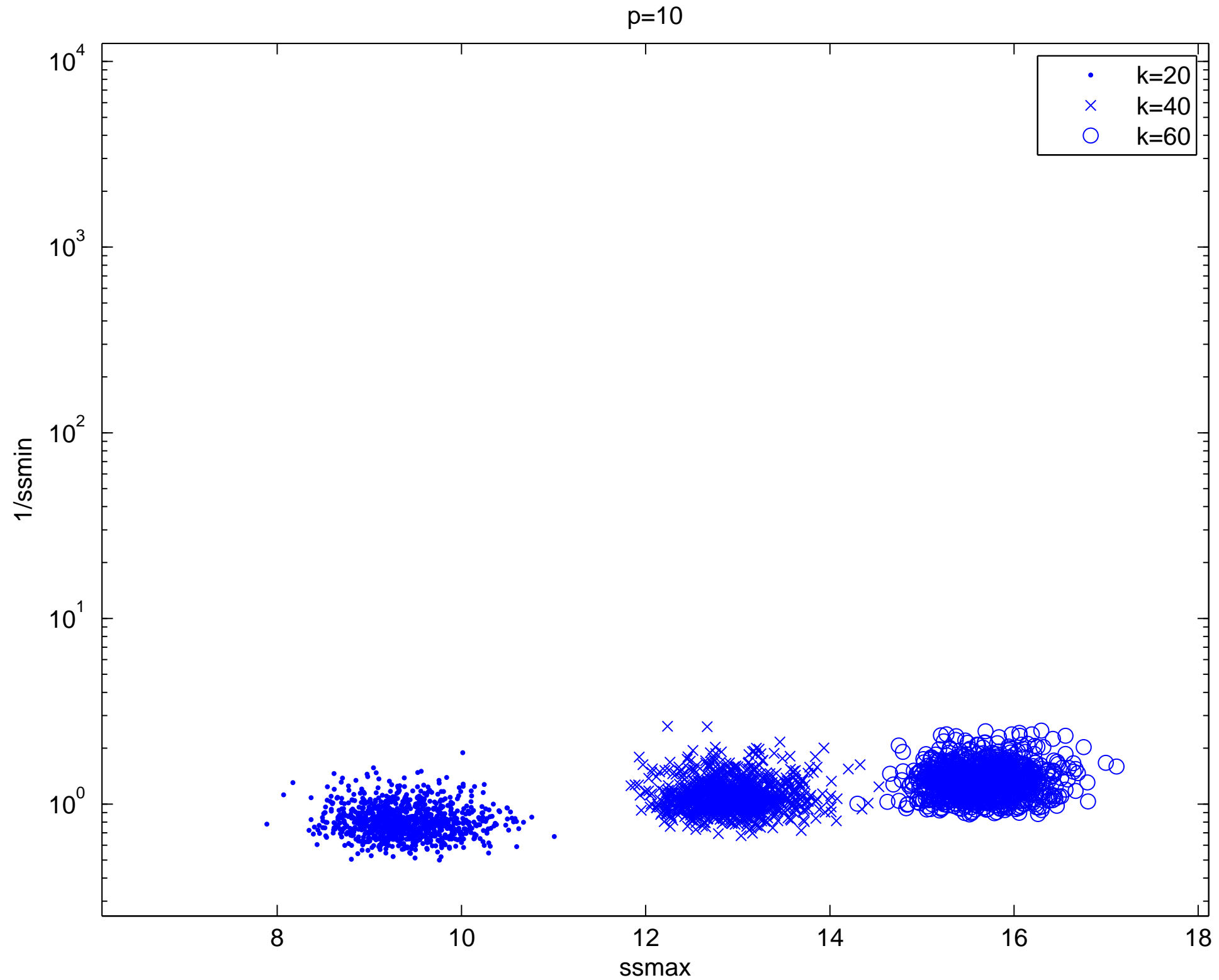
$1/\sigma_{\min}$  is plotted against  $\sigma_{\max}$ .

Scatter plot showing distribution of  $1/\sigma_{\min}$  for  $k \times (k + p)$  Gaussian matrices.  $p = 5$



$1/\sigma_{\min}$  is plotted against  $\sigma_{\max}$ .

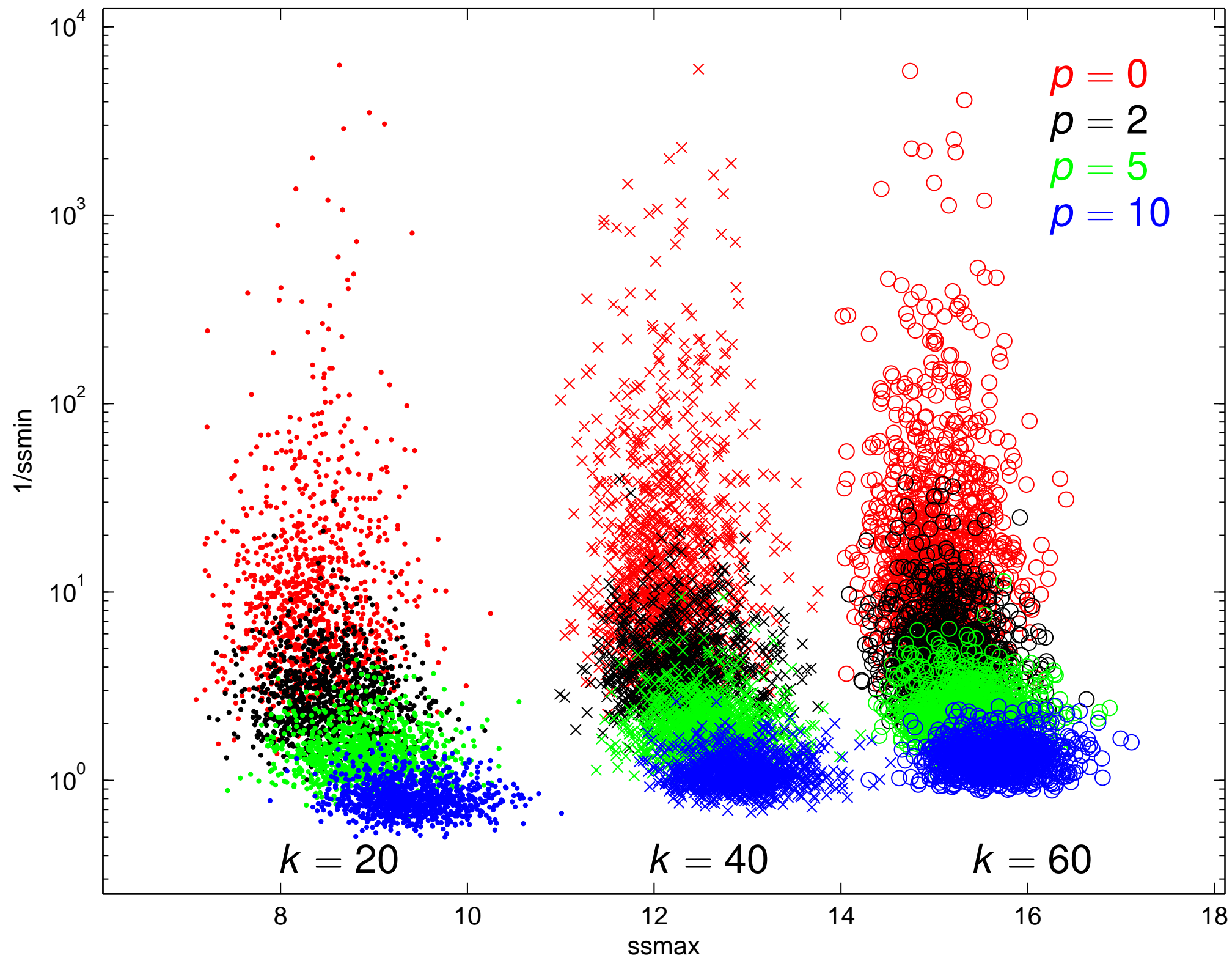
Scatter plot showing distribution of  $1/\sigma_{\min}$  for  $k \times (k + p)$  Gaussian matrices.  $p = 10$



$1/\sigma_{\min}$  is plotted against  $\sigma_{\max}$ .



Scatter plot showing distribution of  $k \times (k + p)$  Gaussian matrices.



$1/\sigma_{\min}$  is plotted against  $\sigma_{\max}$ .

## Simplistic proof that a rectangular Gaussian matrix is well-conditioned:

Let  $\mathbf{G}$  denote a  $k \times \ell$  Gaussian matrix where  $k < \ell$ .

Let “ $g$ ” denote a generic  $\mathcal{N}(0, 1)$  variable and “ $r_j^2$ ” denote a generic  $\chi_j^2$  variable. Then

$$\begin{aligned}
 \mathbf{G} &\sim \begin{bmatrix} g & g & g & g & g & g & \dots \\ g & g & g & g & g & g & \dots \\ g & g & g & g & g & g & \dots \\ g & g & g & g & g & g & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \dots \end{bmatrix} \sim \begin{bmatrix} r_\ell & 0 & 0 & 0 & 0 & 0 & \dots \\ g & g & g & g & g & g & \dots \\ g & g & g & g & g & g & \dots \\ g & g & g & g & g & g & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \dots \end{bmatrix} \\
 &\sim \begin{bmatrix} r_\ell & 0 & 0 & 0 & 0 & 0 & \dots \\ r_{k-1} & g & g & g & g & g & \dots \\ 0 & g & g & g & g & g & \dots \\ 0 & g & g & g & g & g & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \dots \end{bmatrix} \sim \begin{bmatrix} r_\ell & 0 & 0 & 0 & 0 & \dots \\ r_{k-1} & r_{\ell-1} & 0 & 0 & 0 & \dots \\ 0 & g & g & g & g & \dots \\ 0 & g & g & g & g & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \dots \end{bmatrix} \\
 &\sim \begin{bmatrix} r_\ell & 0 & 0 & 0 & 0 & \dots \\ r_{k-1} & r_{\ell-1} & 0 & 0 & 0 & \dots \\ 0 & r_{k-2} & g & g & g & \dots \\ 0 & 0 & g & g & g & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \dots \end{bmatrix} \sim \dots \sim \begin{bmatrix} r_\ell & 0 & 0 & 0 & 0 & \dots \\ r_{k-1} & r_{\ell-1} & 0 & 0 & 0 & \dots \\ 0 & r_{k-2} & r_{\ell-2} & 0 & 0 & \dots \\ 0 & 0 & r_{k-3} & r_{\ell-3} & 0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \dots \end{bmatrix}
 \end{aligned}$$

Gershgorin’s circle theorem will now show that  $\mathbf{G}$  is well-conditioned if, e.g.,  $\ell = 2k$ .

More sophisticated methods are required to get to  $\ell = k + 2$ .

**Some results on Gaussian matrices.** Adapted from HMT 2009/2011; Gordon (1985,1988) for Proposition 1; Chen & Dongarra (2005) for Propositions 2 and 4; Bogdanov (1998) for Proposition 3.

**Proposition 1:** Let  $\mathbf{G}$  be a Gaussian matrix. Then

$$\begin{aligned} (\mathbb{E} [\|\mathbf{S}\mathbf{G}\mathbf{T}\|_{\mathbb{F}}^2])^{1/2} &\leq \|\mathbf{S}\|_{\mathbb{F}} \|\mathbf{T}\|_{\mathbb{F}} \\ \mathbb{E} [\|\mathbf{S}\mathbf{G}\mathbf{T}\|] &\leq \|\mathbf{S}\| \|\mathbf{T}\|_{\mathbb{F}} + \|\mathbf{S}\|_{\mathbb{F}} \|\mathbf{T}\| \end{aligned}$$

**Proposition 2:** Let  $\mathbf{G}$  be a Gaussian matrix of size  $k \times k + p$  where  $p \geq 2$ . Then

$$\begin{aligned} (\mathbb{E} [\|\mathbf{G}^\dagger\|_{\mathbb{F}}^2])^{1/2} &\leq \sqrt{\frac{k}{p-1}} \\ \mathbb{E} [\|\mathbf{G}^\dagger\|] &\leq \frac{e\sqrt{k+p}}{p}. \end{aligned}$$

**Proposition 3:** Suppose  $h$  is Lipschitz  $|h(\mathbf{X}) - h(\mathbf{Y})| \leq L\|\mathbf{X} - \mathbf{Y}\|_{\mathbb{F}}$  and  $\mathbf{G}$  is Gaussian. Then

$$\mathbb{P}[h(\mathbf{G}) > \mathbb{E}[h(\mathbf{G})] + Lu] \leq e^{-u^2/2}.$$

**Proposition 4:** Suppose  $\mathbf{G}$  is Gaussian of size  $k \times k + p$  with  $p \geq 4$ . Then for  $t \geq 1$ :

$$\begin{aligned} \mathbb{P} \left[ \|\mathbf{G}^\dagger\|_{\mathbb{F}} \geq \sqrt{\frac{3k}{p+1}} t \right] &\leq t^{-p} \\ \mathbb{P} \left[ \|\mathbf{G}^\dagger\| \geq \frac{e\sqrt{k+p}}{p+1} t \right] &\leq t^{-(p+1)} \end{aligned}$$

**Recall:**  $\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|^2 \leq \|\mathbf{D}_2\|^2 + \|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\|^2$ , where  $\mathbf{R}_1$  and  $\mathbf{R}_2$  are Gaussian and  $\mathbf{R}_1$  is  $k \times k + p$ .

**Theorem:**  $\mathbb{E}[\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|] \leq \left(1 + \sqrt{\frac{k}{p-1}}\right) \sigma_{k+1} + \frac{e\sqrt{k+p}}{p} \left(\sum_{j=k+1}^{\min(m,n)} \sigma_j^2\right)^{1/2}$ .

**Proof:** First observe that

$$\mathbb{E}\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| = \mathbb{E}(\|\mathbf{D}_2\|^2 + \|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\|^2)^{1/2} \leq \|\mathbf{D}_2\| + \mathbb{E}\|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\|.$$

Condition on  $\mathbf{R}_1$  and use Proposition 1:

$$\begin{aligned} \mathbb{E}\|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\| &\leq \mathbb{E}[\|\mathbf{D}_2\| \|\mathbf{R}_1^\dagger\|_F + \|\mathbf{D}_2\|_F \|\mathbf{R}_1^\dagger\|] \\ &\leq \{\text{H\"older}\} \leq \|\mathbf{D}_2\| (\mathbb{E}\|\mathbf{R}_1^\dagger\|_F^2)^{1/2} + \|\mathbf{D}_2\|_F \mathbb{E}\|\mathbf{R}_1^\dagger\|. \end{aligned}$$

Proposition 2 now provides bounds for  $\mathbb{E}\|\mathbf{R}_1^\dagger\|_F^2$  and  $\mathbb{E}\|\mathbf{R}_1^\dagger\|$  and we get

$$\mathbb{E}\|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\| \leq \sqrt{\frac{k}{p-1}} \|\mathbf{D}_2\| + \frac{e\sqrt{k+p}}{p} \|\mathbf{D}_2\|_F = \sqrt{\frac{k}{p-1}} \sigma_{k+1} + \frac{e\sqrt{k+p}}{p} \left(\sum_{j>k} \sigma_j^2\right)^{1/2}.$$

**Some results on Gaussian matrices.** Adapted from HMT2009/2011; Gordon (1985,1988) for Proposition 1; Chen & Dongarra (2005) for Propositions 2 and 4; Bogdanov (1998) for Proposition 3.

**Proposition 1:** Let  $\mathbf{G}$  be a Gaussian matrix. Then

$$\begin{aligned} (\mathbb{E}[\|\mathbf{S}\mathbf{G}\mathbf{T}\|_{\mathbb{F}}^2])^{1/2} &\leq \|\mathbf{S}\|_{\mathbb{F}} \|\mathbf{T}\|_{\mathbb{F}} \\ \mathbb{E}[\|\mathbf{S}\mathbf{G}\mathbf{T}\|] &\leq \|\mathbf{S}\| \|\mathbf{T}\|_{\mathbb{F}} + \|\mathbf{S}\|_{\mathbb{F}} \|\mathbf{T}\| \end{aligned}$$

**Proposition 2:** Let  $\mathbf{G}$  be a Gaussian matrix of size  $k \times k + p$  where  $p \geq 2$ . Then

$$\begin{aligned} (\mathbb{E}[\|\mathbf{G}^\dagger\|_{\mathbb{F}}^2])^{1/2} &\leq \sqrt{\frac{k}{p-1}} \\ \mathbb{E}[\|\mathbf{G}^\dagger\|] &\leq \frac{e\sqrt{k+p}}{p}. \end{aligned}$$

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$$\mathbb{P}[h(\mathbf{G}) > \mathbb{E}[h(\mathbf{G})] + Lu] \leq e^{-u^2/2}.$$

**Proposition 4:** Suppose  $\mathbf{G}$  is Gaussian of size  $k \times k + p$  with  $p \geq 4$ . Then for  $t \geq 1$ :

$$\begin{aligned} \mathbb{P}\left[\|\mathbf{G}^\dagger\|_{\mathbb{F}} \geq \sqrt{\frac{3k}{p+1}}t\right] &\leq t^{-p} \\ \mathbb{P}\left[\|\mathbf{G}^\dagger\| \geq \frac{e\sqrt{k+p}}{p+1}t\right] &\leq t^{-(p+1)} \end{aligned}$$

**Recall:**  $\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|^2 \leq \|\mathbf{D}_2\|^2 + \|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\|^2$ , where  $\mathbf{R}_1$  and  $\mathbf{R}_2$  are Gaussian and  $\mathbf{R}_1$  is  $k \times k + p$ .

**Theorem:** With probability at least  $1 - 2t^{-p} - e^{-u^2/2}$  it holds that

$$\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\| \leq \left(1 + t \sqrt{\frac{3k}{p+1}} + ut \frac{e\sqrt{k+p}}{p+1}\right) \sigma_{k+1} + \frac{te\sqrt{k+p}}{p+1} \left(\sum_{j>k} \sigma_j^2\right)^{1/2}.$$

**Proof:** Set  $E_t = \left\{ \|\mathbf{R}_1\| \leq \frac{e\sqrt{k+p}}{p+1}t \text{ and } \|\mathbf{R}_1^\dagger\|_F \leq \sqrt{\frac{3k}{p+1}}t \right\}$ . By Proposition 4:  $\mathbb{P}(E_t^c) \leq 2t^{-p}$ .

Set  $h(\mathbf{X}) = \|\mathbf{D}_2\mathbf{X}\mathbf{R}_1^\dagger\|$ . A direct calculation shows

$$|h(\mathbf{X}) - h(\mathbf{Y})| \leq \|\mathbf{D}_2\| \|\mathbf{R}_1^\dagger\| \|\mathbf{X} - \mathbf{Y}\|_F.$$

Hold  $\mathbf{R}_1$  fixed and take the expectation on  $\mathbf{R}_2$ . Then Proposition 1 applies and so

$$\mathbb{E}[h(\mathbf{R}_2) \mid \mathbf{R}_1] \leq \|\mathbf{D}_2\| \|\mathbf{R}_1^\dagger\|_F + \|\mathbf{D}_2\|_F \|\mathbf{R}_1^\dagger\|.$$

Now use Proposition 3 (concentration of measure)

$$\mathbb{P}\left[\underbrace{\|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\|}_{=h(\mathbf{R}_2)} > \underbrace{\|\mathbf{D}_2\| \|\mathbf{R}_1^\dagger\|_F + \|\mathbf{D}_2\|_F \|\mathbf{R}_1^\dagger\|}_{=\mathbb{E}[h(\mathbf{R}_2)]} + \underbrace{\|\mathbf{D}_2\| \|\mathbf{R}_1^\dagger\|}_{=L} u \mid E_t\right] < e^{-u^2/2}.$$

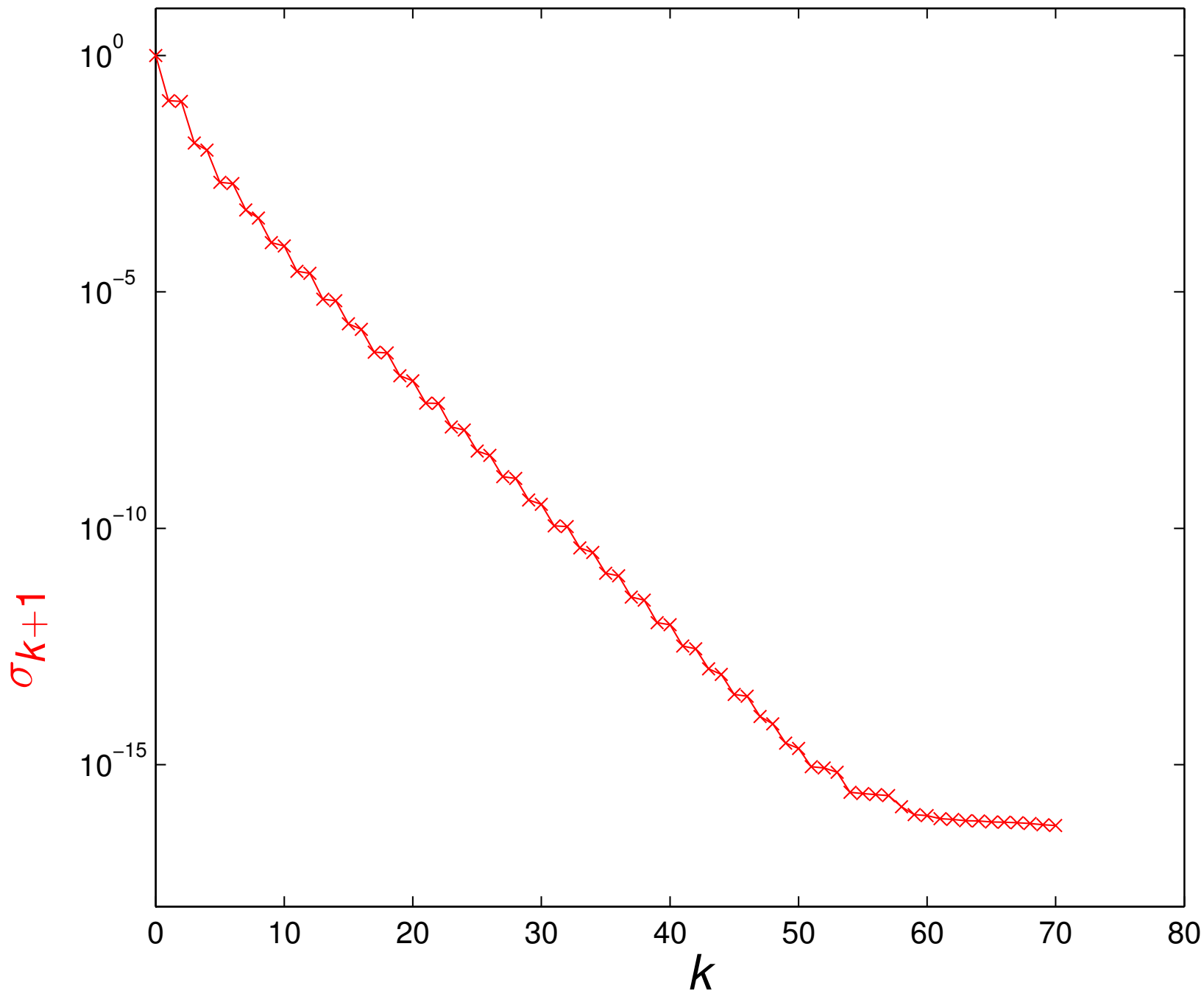
When  $E_t$  holds true, we have bounds on the “badness” of  $\mathbf{R}_1^\dagger$ :

$$\mathbb{P}\left[\|\mathbf{D}_2\mathbf{R}_2\mathbf{R}_1^\dagger\| > \|\mathbf{D}_2\| \sqrt{\frac{3k}{p+1}}t + \|\mathbf{D}_2\|_F \frac{e\sqrt{k+p}}{p+1}t + \|\mathbf{D}_2\| \frac{e\sqrt{k+p}}{p+1}ut \mid E_t\right] < e^{-u^2/2}.$$

The theorem is obtained by using  $\mathbb{P}(E_t^c) \leq 2t^{-p}$  to remove the conditioning of  $E_t$ .

## Example 1:

We consider a  $1\,000 \times 1\,000$  matrix  $\mathbf{A}$  whose singular values are shown below:



The red line indicates the singular values  $\sigma_{k+1}$  of  $\mathbf{A}$ . These indicate the theoretically minimal approximation error.

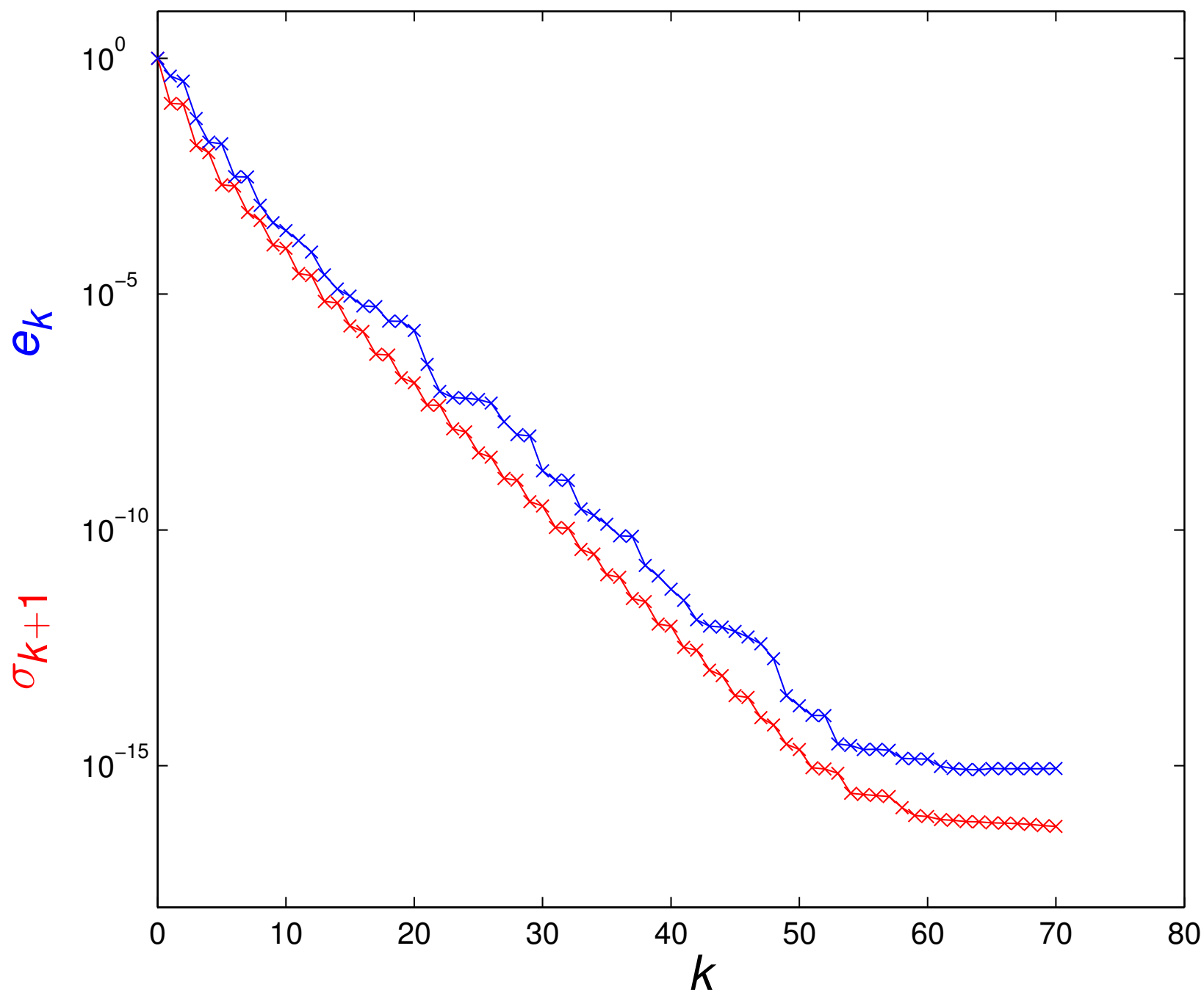
---

$\mathbf{A}$  is a discrete approximation of a certain compact integral operator normalized so that  $\|\mathbf{A}\| = 1$ .

Curiously, the nature of  $\mathbf{A}$  is in a strong sense irrelevant: the error distribution depends only on  $\{\sigma_j\}_{j=1}^{\min(m,n)}$ .

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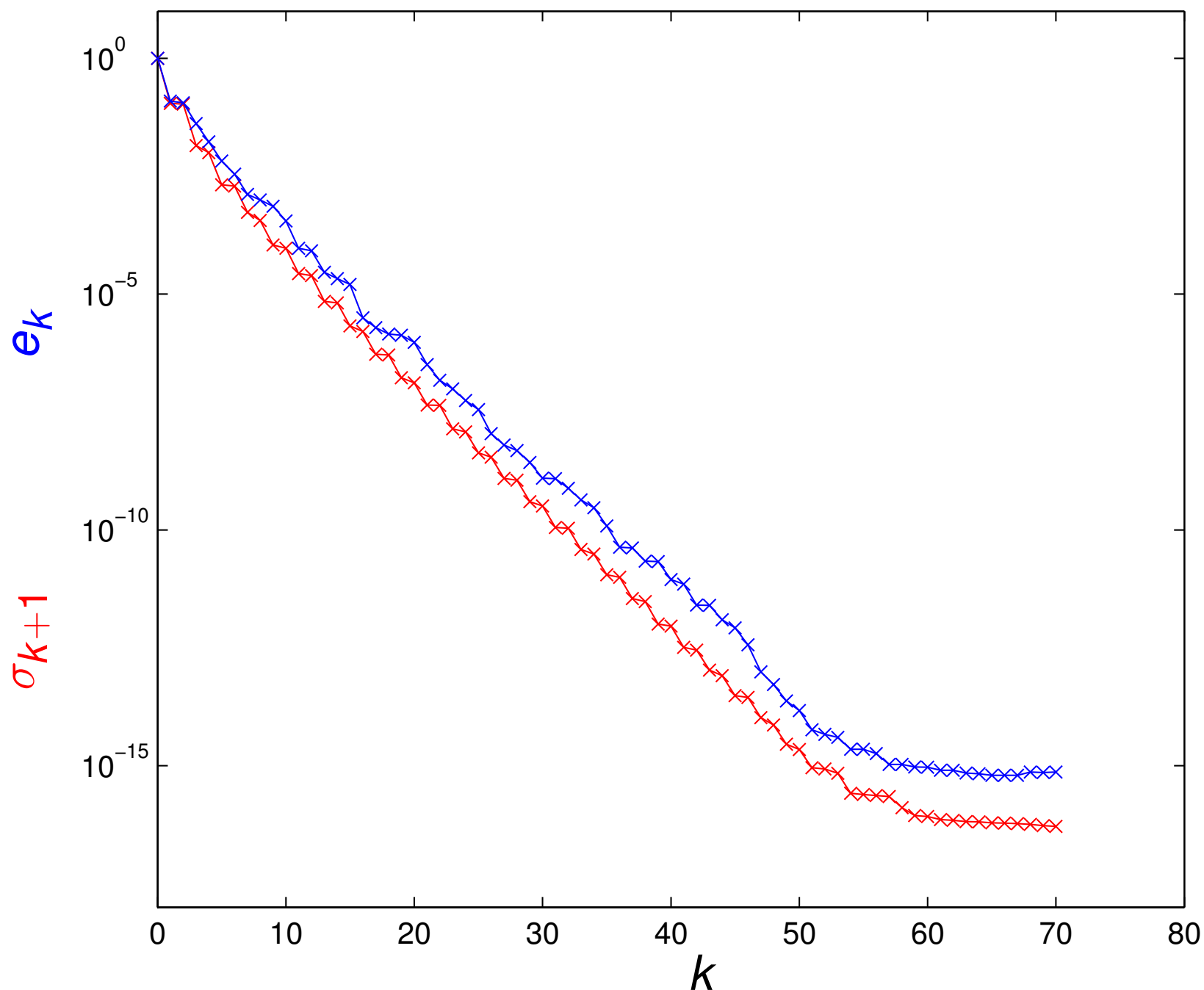
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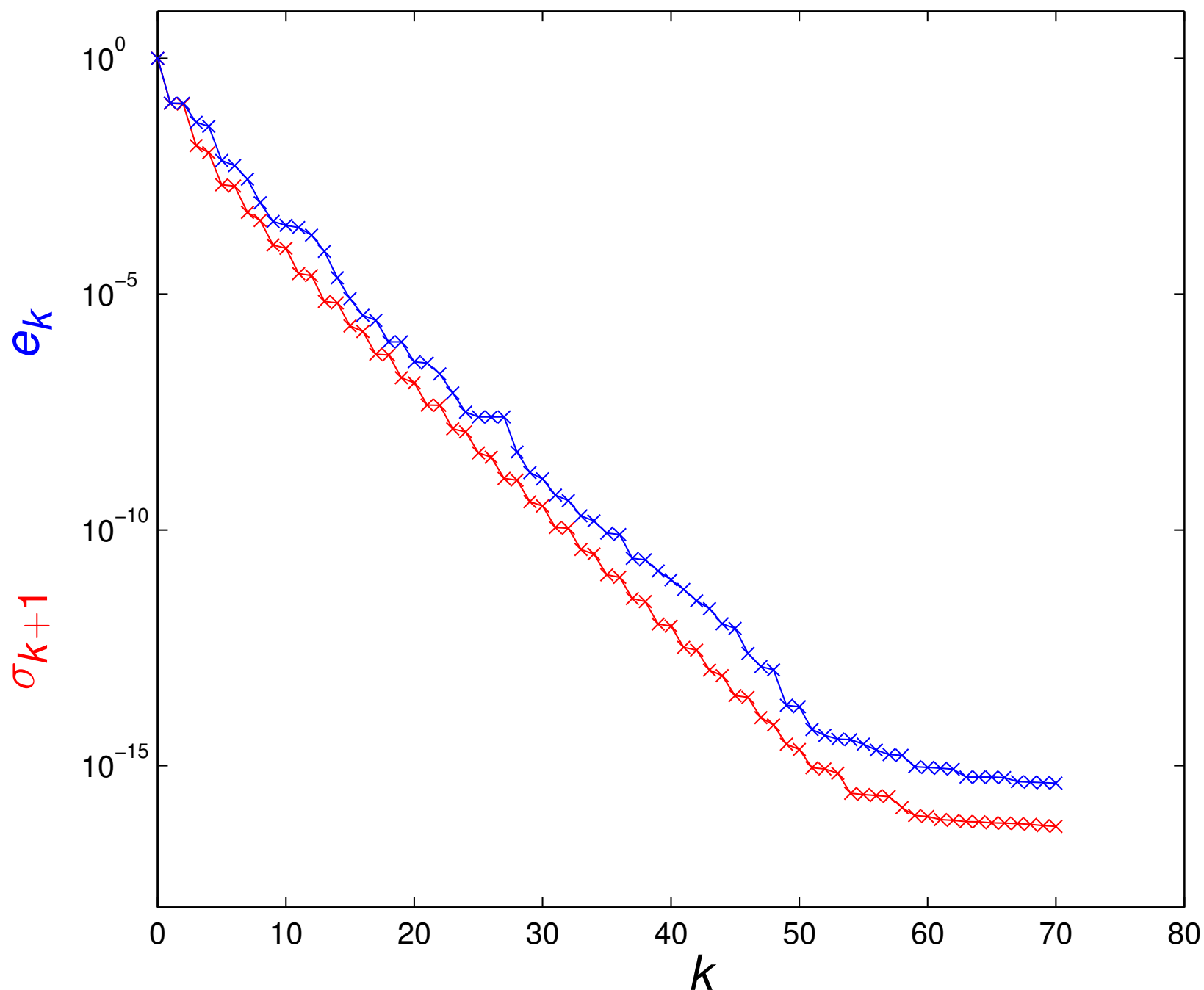
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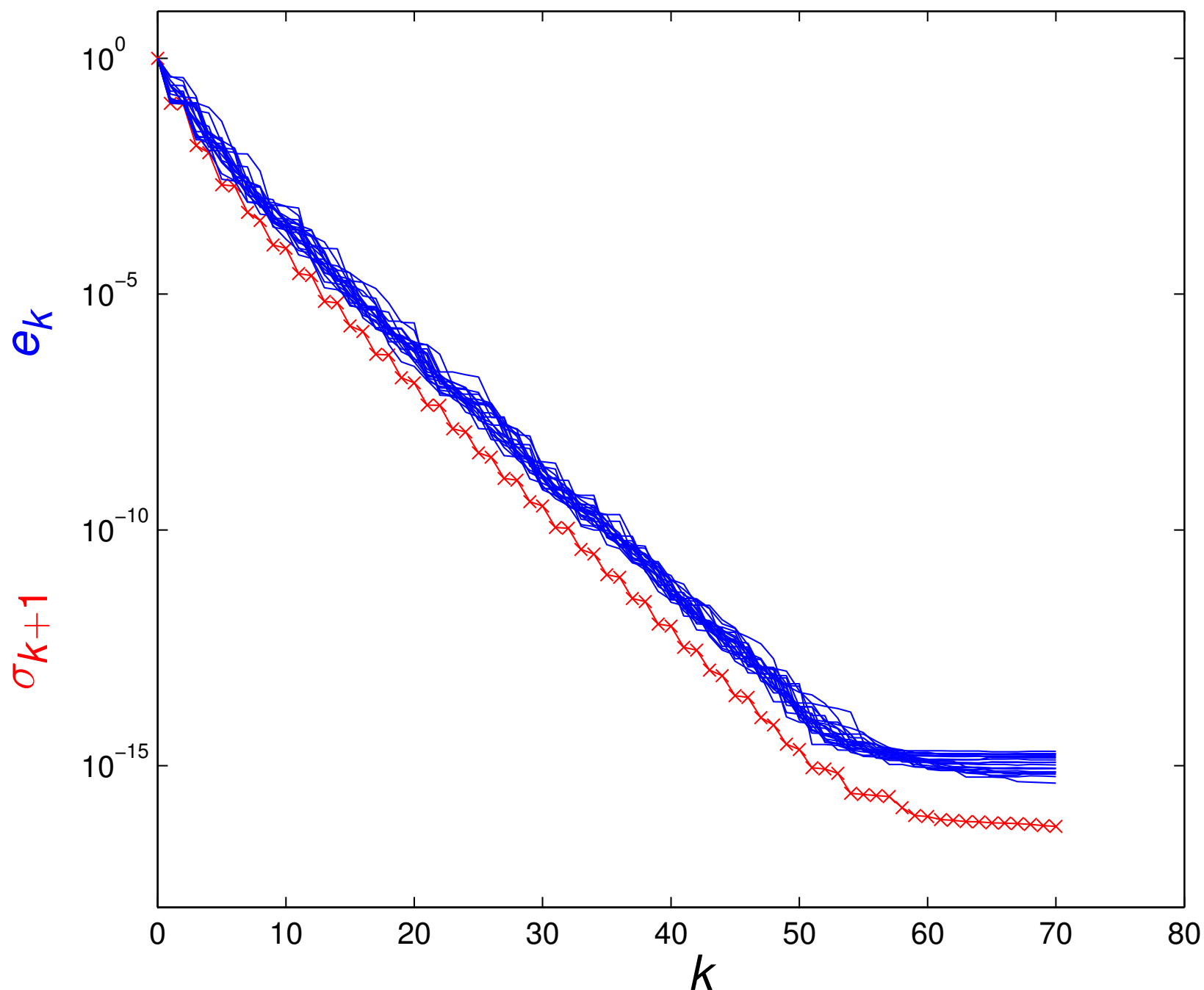
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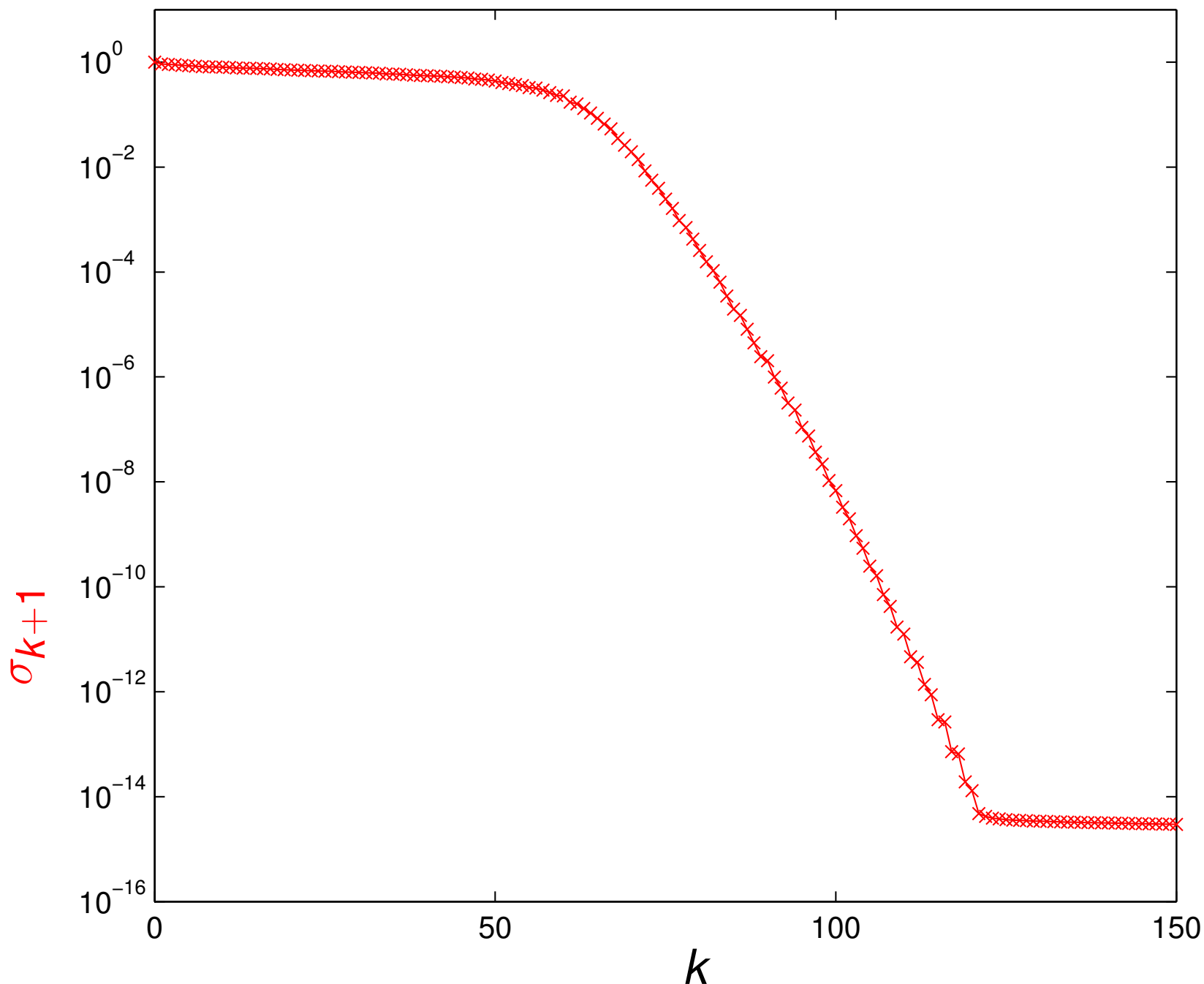
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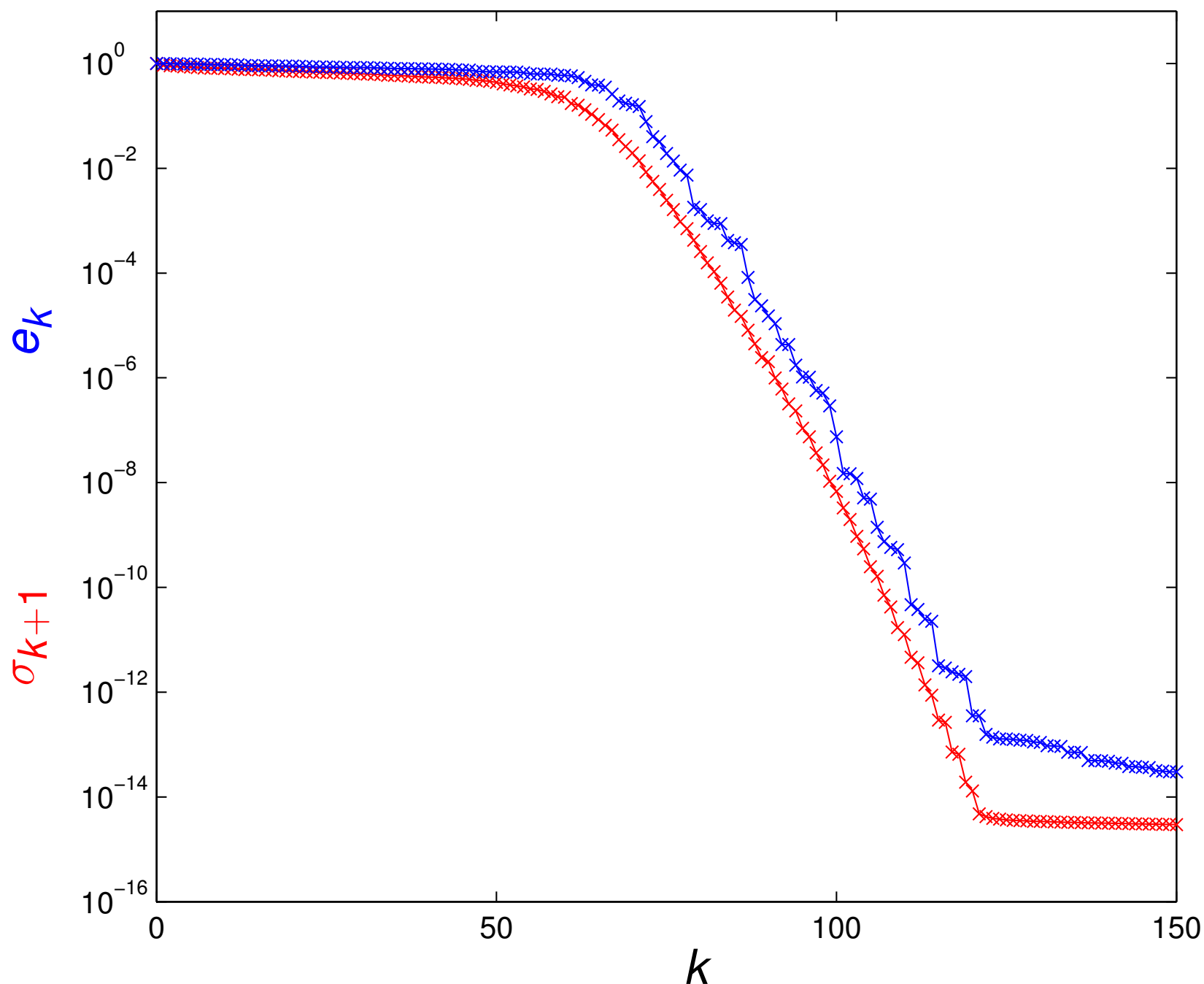
---

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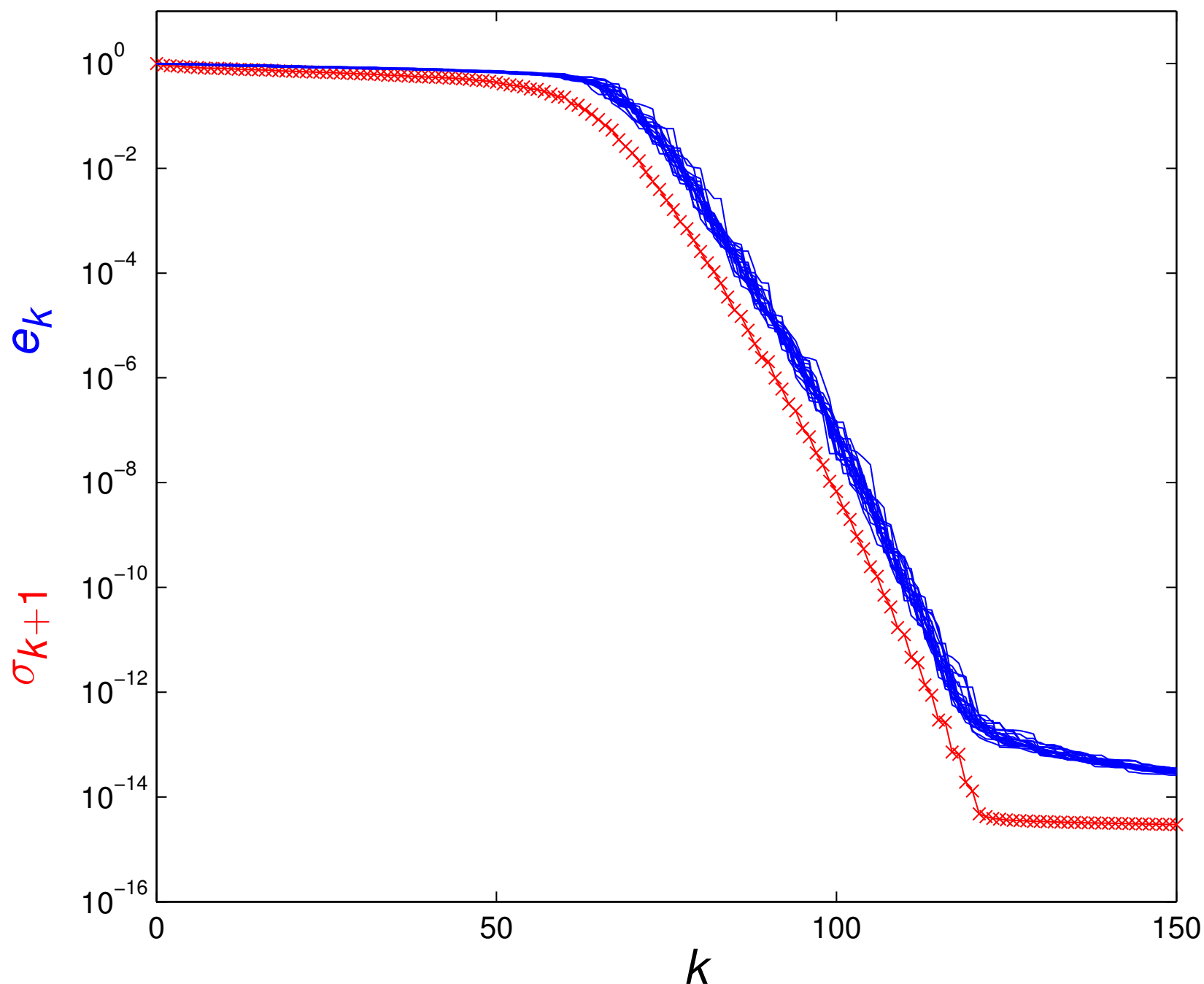
The blue line indicates the actual errors  $e_k$  incurred by one instantiation of the proposed method.

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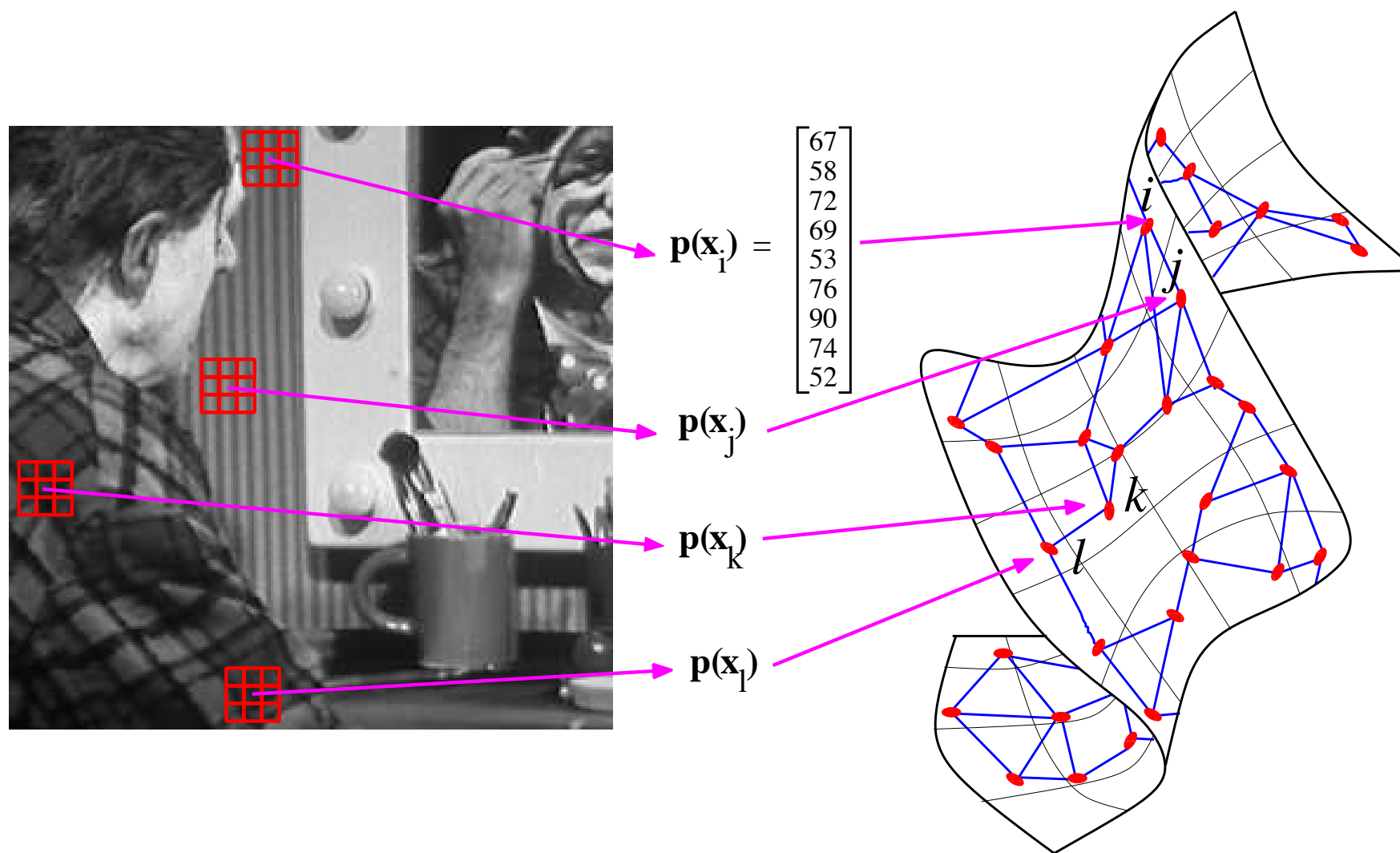
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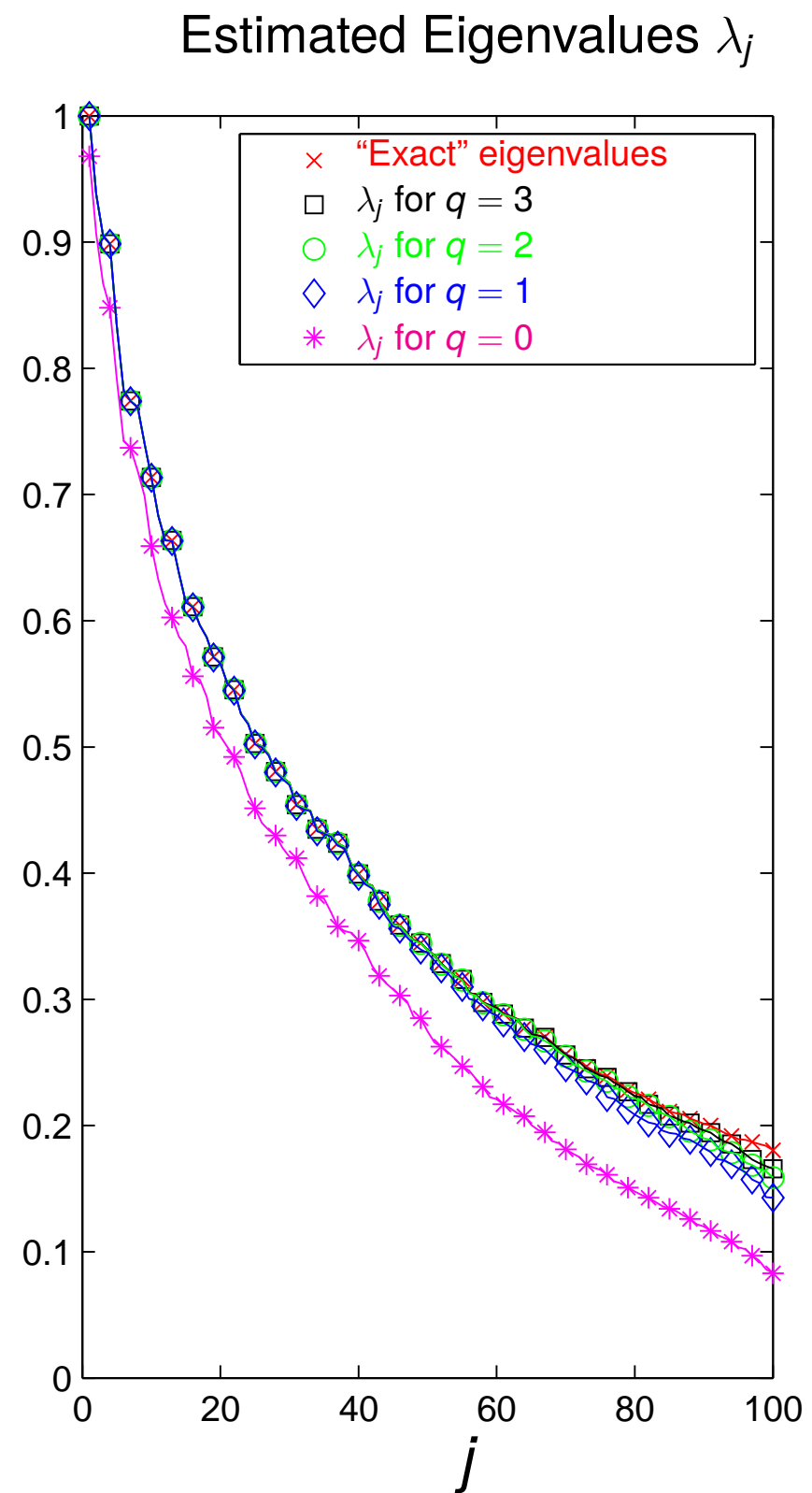
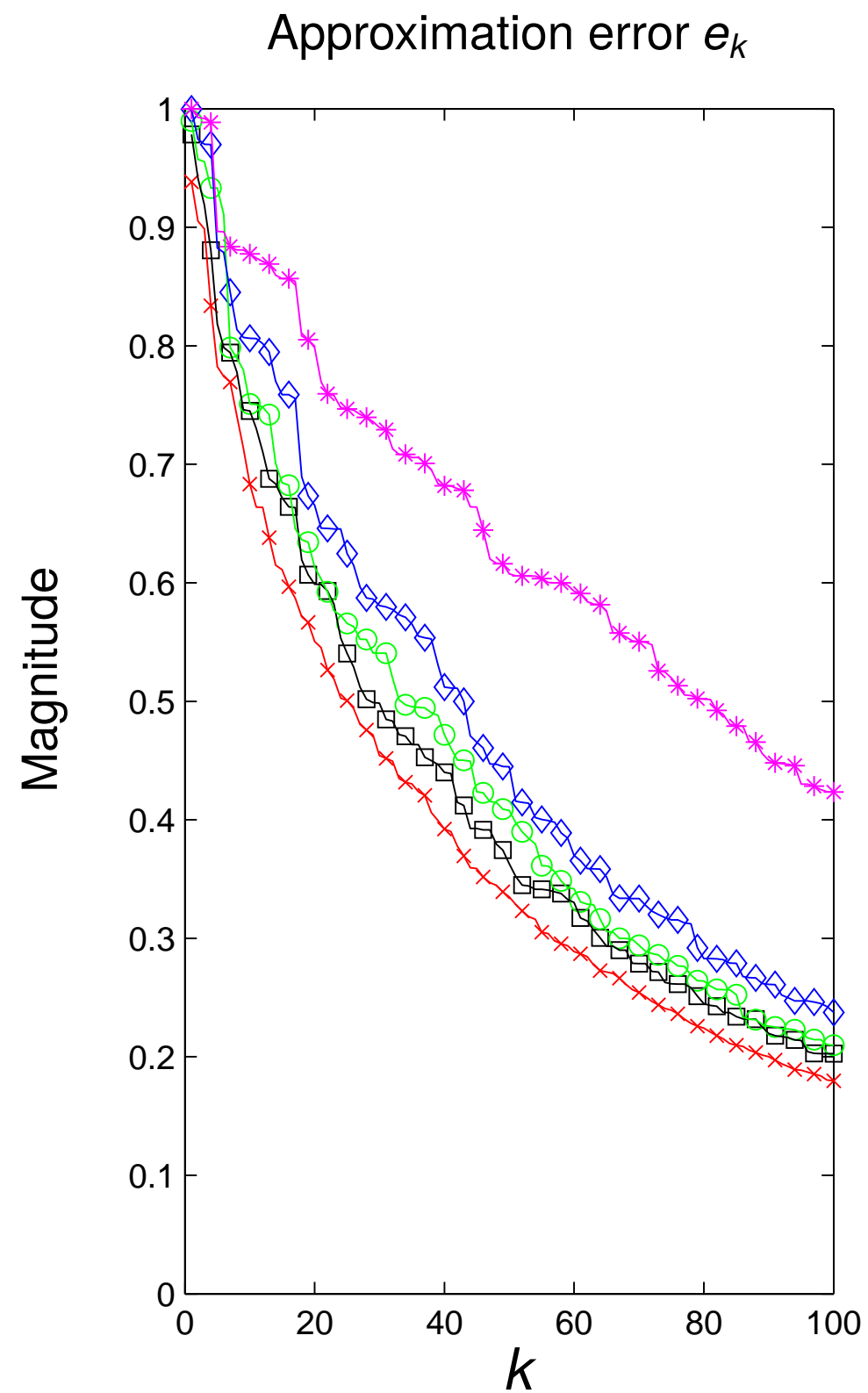
### Example 3:

The matrix  $\mathbf{A}$  being analyzed is a  $9025 \times 9025$  matrix arising in a diffusion geometry approach to image processing.

To be precise,  $\mathbf{A}$  is a graph Laplacian on the manifold of  $3 \times 3$  patches.



*Joint work with François Meyer of the University of Colorado at Boulder.*



The pink lines illustrates the performance of the basic random sampling scheme. The errors are huge, and the estimated eigenvalues are much too small.



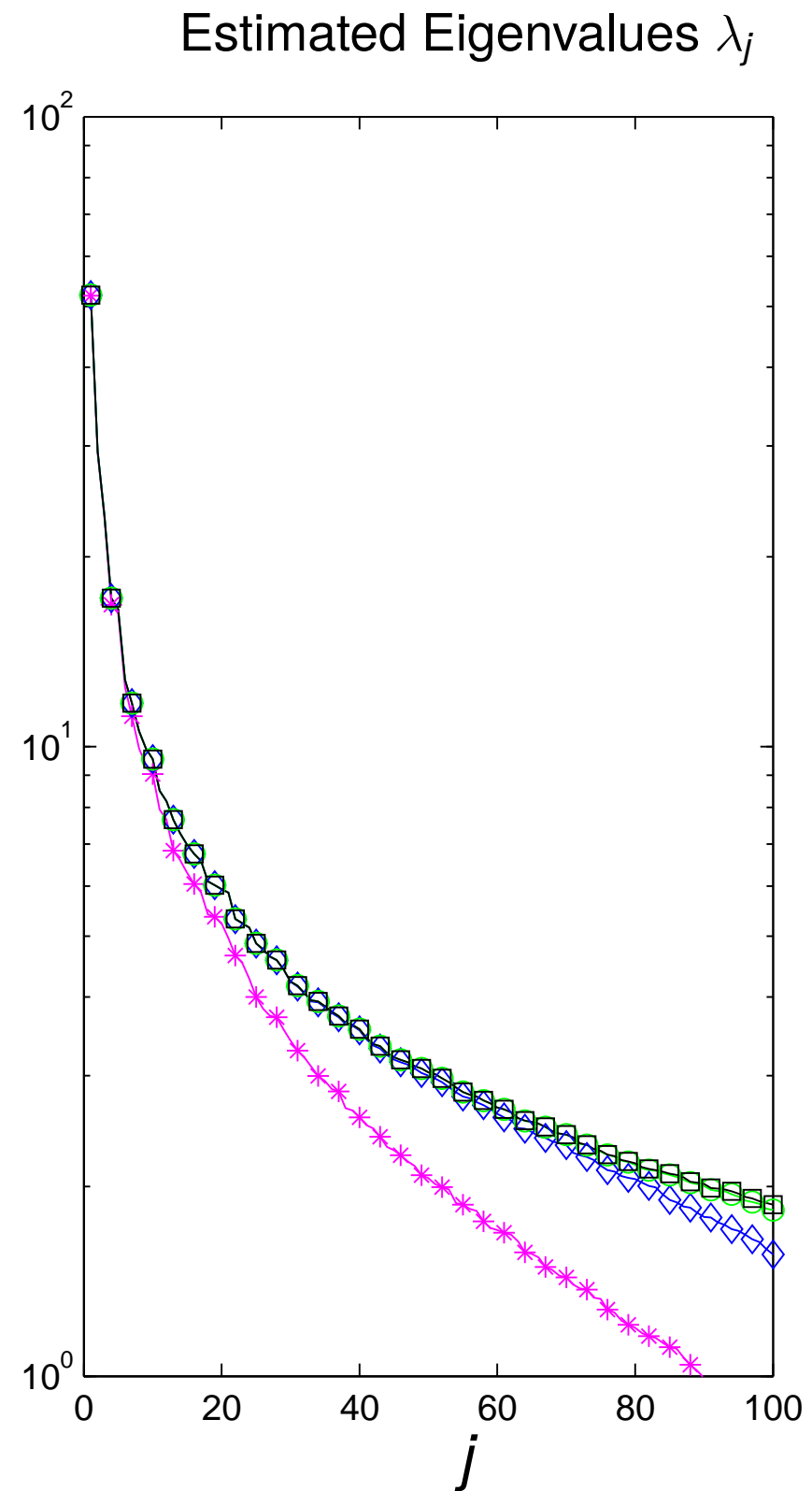
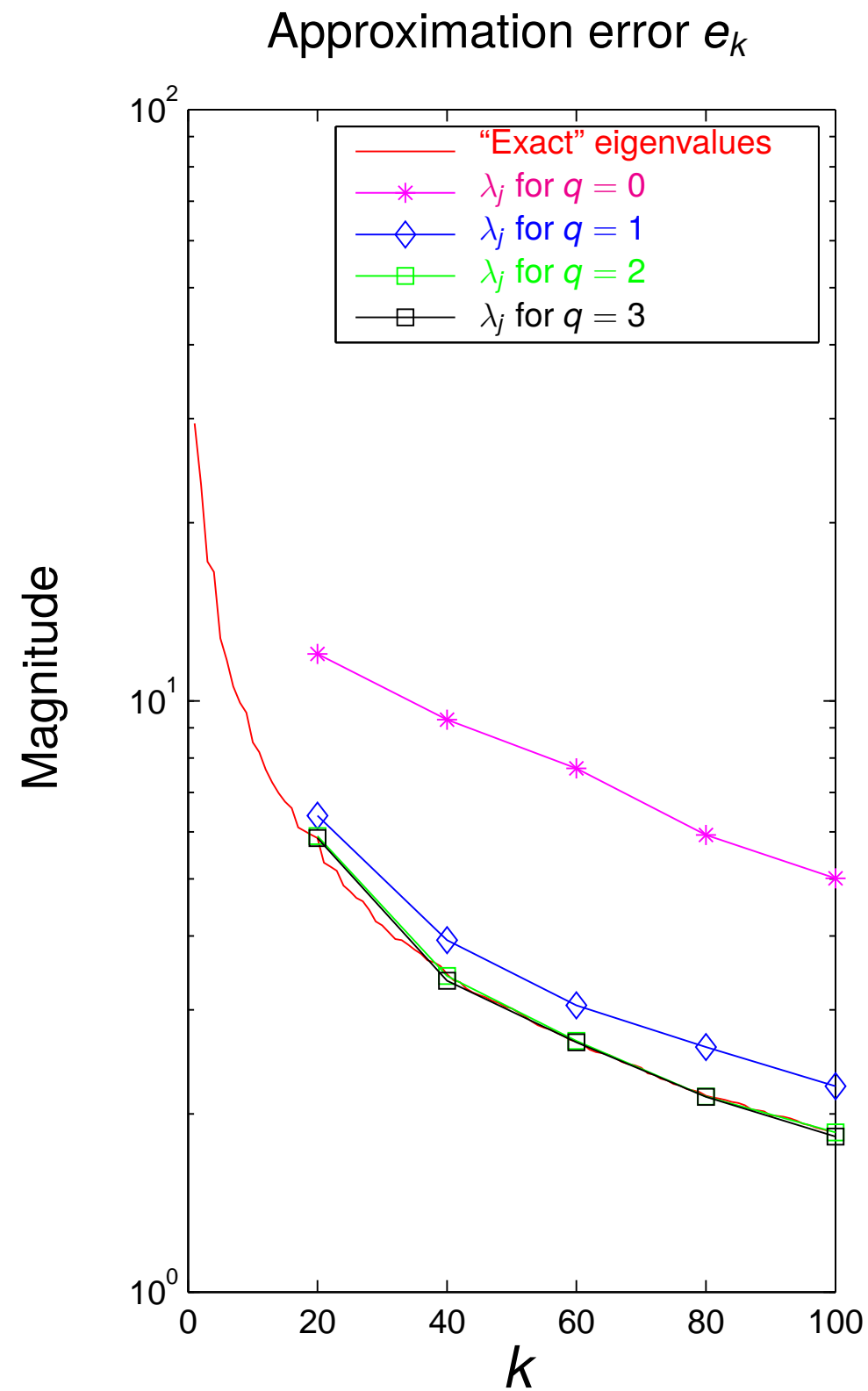
## Example 4: “Eigenfaces”

We next process a data base containing  $m = 7\,254$  pictures of faces

Each image consists of  $n = 384 \times 256 = 98\,304$  gray scale pixels.

We center and scale the pixels in each image, and let the resulting values form a column of a  $98\,304 \times 7\,254$  data matrix  $\mathbf{A}$ .

The left singular vectors of  $\mathbf{A}$  are the so called *eigenfaces* of the data base.



The pink lines illustrates the performance of the basic random sampling scheme. Again, the errors are huge, and the estimated eigenvalues are much too small.

## Power method for improving accuracy:

The error depends on how quickly the singular values decay. Recall that

$$\mathbb{E} \|\mathbf{A} - \mathbf{Q}\mathbf{Q}^* \mathbf{A}\| \leq \left(1 + \sqrt{\frac{k}{p-1}}\right) \sigma_{k+1} + \frac{e \sqrt{k+p}}{p} \left( \sum_{j=k+1}^{\min(m,n)} \sigma_j^2 \right)^{1/2}.$$

The faster the singular values decay — the stronger the relative weight of the dominant modes in the samples.

**Idea:** The matrix  $(\mathbf{A}\mathbf{A}^*)^q \mathbf{A}$  has the same left singular vectors as  $\mathbf{A}$ , and singular values

$$\sigma_j((\mathbf{A}\mathbf{A}^*)^q \mathbf{A}) = (\sigma_j(\mathbf{A}))^{2q+1}.$$

Much faster decay — so let us use the sample matrix

$$\mathbf{Y} = (\mathbf{A}\mathbf{A}^*)^q \mathbf{A} \mathbf{R}$$

instead of

$$\mathbf{Y} = \mathbf{A} \mathbf{R}.$$

**References:** Paper by Rokhlin, Szlam, Tygert (2008). Suggestions by Ming Gu. Also similar to “block power method,” “block Lanczos,” “subspace iteration.”

*Input:* An  $m \times n$  matrix  $\mathbf{A}$ , a target rank  $\ell$ , and a small integer  $q$ .

*Output:* Rank- $\ell$  factors  $\mathbf{U}$ ,  $\mathbf{D}$ , and  $\mathbf{V}$  in an approximate SVD  $\mathbf{A} \approx \mathbf{UDV}^*$ .

(1) Draw an  $n \times \ell$  **random matrix**  $\mathbf{R}$ .

(2) Form the  $m \times \ell$  **sample matrix**  $\mathbf{Y} = (\mathbf{A}\mathbf{A}^*)^q \mathbf{A}\mathbf{R}$ .

(3) Compute an **ON matrix**  $\mathbf{Q}$  s.t.  $\mathbf{Y} = \mathbf{Q}\mathbf{Q}^*\mathbf{Y}$ .

(4) Form the small matrix  $\mathbf{B} = \mathbf{Q}^* \mathbf{A}$ .

(5) Factor the small matrix  $\mathbf{B} = \hat{\mathbf{U}}\mathbf{D}\mathbf{V}^*$ .

(6) Form  $\mathbf{U} = \mathbf{Q}\hat{\mathbf{U}}$ .

- Detailed (and, we believe, close to sharp) error bounds have been proven.

For instance, with  $\mathbf{A}^{\text{computed}} = \mathbf{UDV}^*$ , the expectation of the error satisfies:

$$(1) \quad \mathbb{E} \left[ \|\mathbf{A} - \mathbf{A}^{\text{computed}}\| \right] \leq \left( 1 + 4\sqrt{\frac{2 \min(m, n)}{k-1}} \right)^{1/(2q+1)} \sigma_{k+1}(\mathbf{A}).$$

*Reference: Halko, Martinsson, Tropp (2011).*

- The improved accuracy from the modified scheme comes at a cost;

$2q + 1$  passes over the matrix are required instead of 1.

However,  $q$  can often be chosen quite small in practice,  $q = 2$  or  $q = 3$ , say.

- The bound (1) assumes exact arithmetic.

To handle round-off errors, variations of subspace iterations can be used.

These are entirely numerically stable and achieve the same error bound.

## A numerically stable version of the “power method”:

*Input:* An  $m \times n$  matrix  $\mathbf{A}$ , a target rank  $\ell$ , and a small integer  $q$ .

*Output:* Rank- $\ell$  factors  $\mathbf{U}$ ,  $\mathbf{D}$ , and  $\mathbf{V}$  in an approximate SVD  $\mathbf{A} \approx \mathbf{UDV}^*$ .

Draw an  $n \times \ell$  Gaussian random matrix  $\mathbf{R}$ .

Set  $\mathbf{Q} = \text{orth}(\mathbf{AR})$

**for**  $i = 1, 2, \dots, q$

$\mathbf{W} = \text{orth}(\mathbf{A}^* \mathbf{Q})$

$\mathbf{Q} = \text{orth}(\mathbf{AW})$

**end for**

$\mathbf{B} = \mathbf{Q}^* \mathbf{A}$

$[\hat{\mathbf{U}}, \mathbf{D}, \mathbf{V}] = \text{svd}(\mathbf{B})$

$\mathbf{U} = \mathbf{Q}\hat{\mathbf{U}}$ .

**Note:** Algebraically, the method with orthogonalizations is identical to the “original” method where  $\mathbf{Q} = \text{orth}((\mathbf{AA}^*)^q \mathbf{AR})$ .

**Note:** This is a classic subspace iteration.

The novelty is the error analysis, and the finding that using a very small  $q$  is often fine.

(In fact, our analysis allows  $q$  to be zero...)

## ADAPTIVE RANK DETERMINATION

How to proceed when the rank of a matrix is not known in advance.

## Adaptive rank determination — vector-by-vector technique

Let us again start by considering the simplistic case where  $\mathbf{A}$  is *exactly* rank-deficient.

Let  $\mathbf{A}$  be an  $m \times n$  matrix of exact rank  $k$ , where  $k$  is *unknown*.

We seek an  $m \times k$  matrix  $\mathbf{Q}$  whose columns form an ON basis for  $\text{col}(\mathbf{A})$ .

$\mathbf{Q} = [ ];$

**for**  $i = 1, 2, 3, \dots, ???$

Draw an  $n \times 1$  Gaussian random vector  $\mathbf{r}_i$ .

Compute an  $m \times 1$  sample vector  $\mathbf{y}_i = \mathbf{A}\mathbf{r}_i$ .

Project the sample vector away from the basis computed  $\mathbf{z}_i = \mathbf{y}_i - \mathbf{Q}\mathbf{Q}^*\mathbf{y}_i$ .

Add the new element to the basis  $\mathbf{Q} = \left[ \mathbf{Q} \frac{\mathbf{z}_i}{\|\mathbf{z}_i\|} \right]$ .

**end for**

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**Observation 1:** While  $i \leq k$ , we know that  $\mathbf{z}_i \neq \mathbf{0}$  with probability 1.



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```

```
end for
```

**Observation 1:** While  $i \leq k$ , we know that  $\mathbf{z}_i \neq \mathbf{0}$  with probability 1.

**Observation 2:** Once you come to step  $i = k + 1$ , the vector  $\mathbf{z}_{k+1}$  must be zero!

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```

```
    Project the sample vector away from the basis computed  $\mathbf{z}_i = \mathbf{y}_i - \mathbf{Q}\mathbf{Q}^*\mathbf{y}_i$ .
```

```
    if  $[\mathbf{z}_i = \mathbf{0}]$  then
```

```
        The rank is  $k = i - 1$ .
```

```
        break
```

```
    else
```

```
        Add the new element to the basis  $\mathbf{Q} = [\mathbf{Q} \frac{\mathbf{z}_i}{\|\mathbf{z}_i\|}]$ .
```

```
    end if
```

```
end for
```

## Adaptive rank determination — vector-by-vector technique

Let  $\mathbf{A}$  be an  $m \times n$  matrix whose singular values decay, but do not hit zero.

Let  $\varepsilon > 0$  be a given tolerance. We seek an  $m \times k$  ON matrix  $\mathbf{Q}$  s.t.  $\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|_{\text{Fro}} \leq \varepsilon$ .

$\mathbf{Q} = [ ]$ ;

**for**  $i = 1, 2, 3, \dots, ???$

Draw an  $n \times 1$  Gaussian random vector  $\mathbf{r}_i$ .

Compute an  $m \times 1$  sample vector  $\mathbf{y}_i = \mathbf{A}\mathbf{r}_i$ .

Project the sample vector away from the basis computed  $\mathbf{z}_i = \mathbf{y}_i - \mathbf{Q}\mathbf{Q}^*\mathbf{y}_i$ .

Add the new element to the basis  $\mathbf{Q} = \left[ \mathbf{Q} \frac{\mathbf{z}_i}{\|\mathbf{z}_i\|} \right]$ .

**end for**

## Adaptive rank determination — vector-by-vector technique

Let  $\mathbf{A}$  be an  $m \times n$  matrix whose singular values decay, but do not hit zero.

Let  $\varepsilon > 0$  be a given tolerance. We seek an  $m \times k$  ON matrix  $\mathbf{Q}$  s.t.  $\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|_{\text{Fro}} \leq \varepsilon$ .

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**end for**

Observe that

$$\mathbf{z}_i = \mathbf{y}_i - \mathbf{Q}\mathbf{Q}^*\mathbf{y}_i = \mathbf{A}\mathbf{r}_i - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\mathbf{r}_i = (\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A})\mathbf{r}_i.$$

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In consequence, since  $\mathbf{r}_i$  is Gaussian,

$$\mathbb{E}[\|\mathbf{z}_i\|^2] = \|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|_{\text{Fro}}^2.$$

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**Observation 1:** Once you observe several consecutive  $\mathbf{z}_i$  such that, say,  $\|\mathbf{z}_i\| \leq \varepsilon/2$ , it will “likely” be the case that  $\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|_{\text{Fro}} \leq \varepsilon$ .

**Observation 2:** You need to **block** the algorithm for computational efficiency.

## Adaptive rank determination

Let  $\mathbf{A}$  be an  $m \times n$  matrix whose singular values decay, but do not hit zero.

Let  $\varepsilon > 0$  be a given tolerance, and let  $b$  be a “block size.”

We seek an  $m \times k$  ON matrix  $\mathbf{Q}$  s.t.  $\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|_{\text{Fro}} \leq \varepsilon$ .

```
Q = [ ];
```

```
for  $i = 1, 2, 3, \dots$ 
```

```
    Draw an  $n \times b$  Gaussian random matrix  $\mathbf{R}_i$ .
```

```
    Compute an  $m \times b$  sample matrix  $\mathbf{Y}_i = \mathbf{A}\mathbf{R}_i$ .
```

```
    Project the sample columns away from the basis computed  $\mathbf{Z}_i = \mathbf{Y}_i - \mathbf{Q}\mathbf{Q}^*\mathbf{Y}_i$ .
```

```
    Orthonormalize the samples  $[\mathbf{Q}_i, \mathbf{R}_i] = \text{qr}(\mathbf{Z}_i, 0)$ . (Unpivoted QR factorization!)
```

```
    if [“several consecutive columns of  $\mathbf{R}_i$  are small”] then
```

```
        Add the appropriate number of columns of  $\mathbf{Q}_i$  to  $\mathbf{Q}$ .
```

```
        break
```

```
    else
```

```
        Add the new element to the basis  $\mathbf{Q} = [\mathbf{Q} \ \mathbf{Q}_i]$ .
```

```
    end if
```

```
end for
```

*Warning: Re-orthogonalization is often needed to combat floating point errors.*

## Adaptive rank determination — with updating

Consider the special case that  $\mathbf{A}$  can be updated, e.g. if it is dense and stored in RAM.

Let  $\mathbf{A}$  be an  $m \times n$  matrix whose singular values decay, but do not hit zero.

Let  $\varepsilon > 0$  be a given tolerance, and let  $b$  be a “block size.”

We seek an  $m \times k$  ON matrix  $\mathbf{Q}$  s.t.  $\|\mathbf{A} - \mathbf{Q}\mathbf{Q}^*\mathbf{A}\|_{\text{Fro}} \leq \varepsilon$ .

```
(1)  $\mathbf{Q} = []$ ;  $\mathbf{B} = []$ ;  
(2) while  $\|\mathbf{A}\| > \varepsilon$   
(3)     Draw an  $n \times b$  Gaussian matrix  $\mathbf{R}_j$ .  
(4)     Compute the  $m \times b$  matrix  $[\mathbf{Q}_j, \sim] = \text{qr}(\mathbf{A}\mathbf{R}_j, 0)$ .  
(5)      $\mathbf{B}_j = \mathbf{Q}_j^*\mathbf{A}$   
(6)      $\mathbf{Q} = [\mathbf{Q} \ \mathbf{Q}_j]$   
(7)      $\mathbf{B} = \begin{bmatrix} \mathbf{B} \\ \mathbf{B}_j \end{bmatrix}$   
(8)      $\mathbf{A} = \mathbf{A} - \mathbf{Q}_j\mathbf{B}_j$   
(9) end while
```

A blocked and randomized variation of the classical “modified Gram-Schmidt” algorithm.

*Warning: Re-orthogonalization is often needed to combat floating point errors.*



## Adaptive rank determination — with updating

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(9) end while
```

**Observation:** Almost all the work is done by matrix-matrix multiplies.

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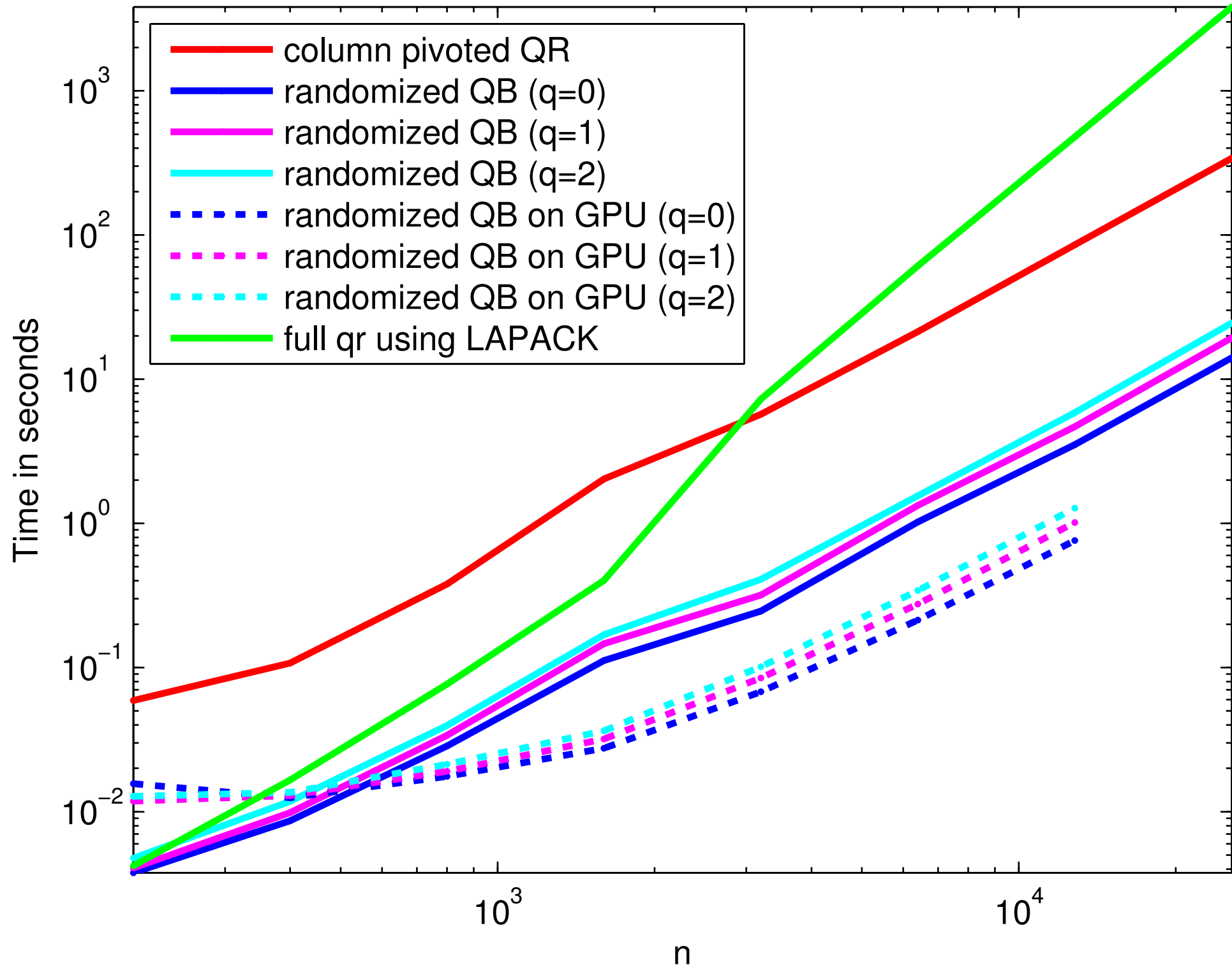
```
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**Observation:** Almost all the work is done by matrix-matrix multiplies.



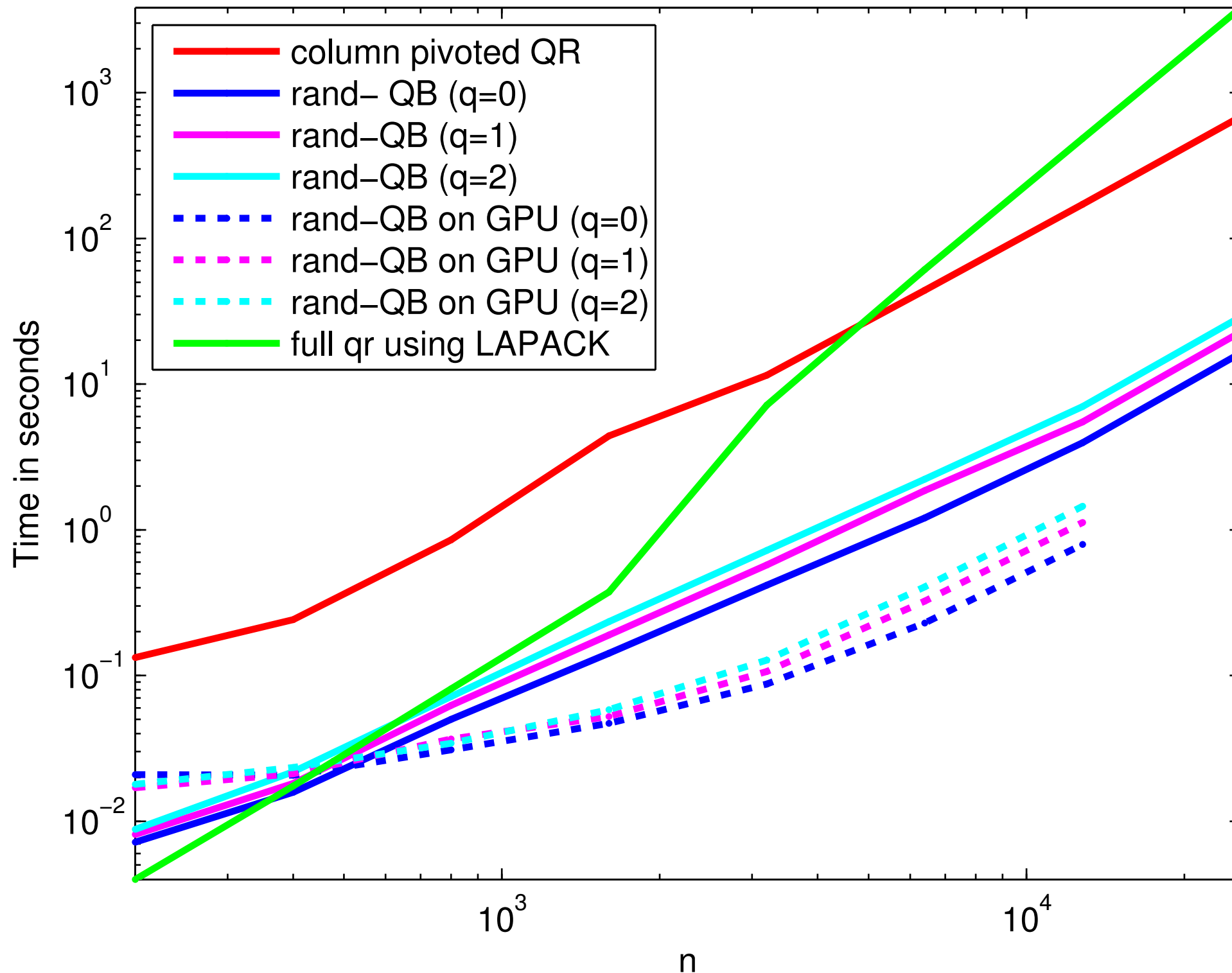
*This algorithm is ideal for running on modern CPUs and GPUs!*

Time for compression of  $n \times n$  matrix.  $k=100$   $kstep=20$



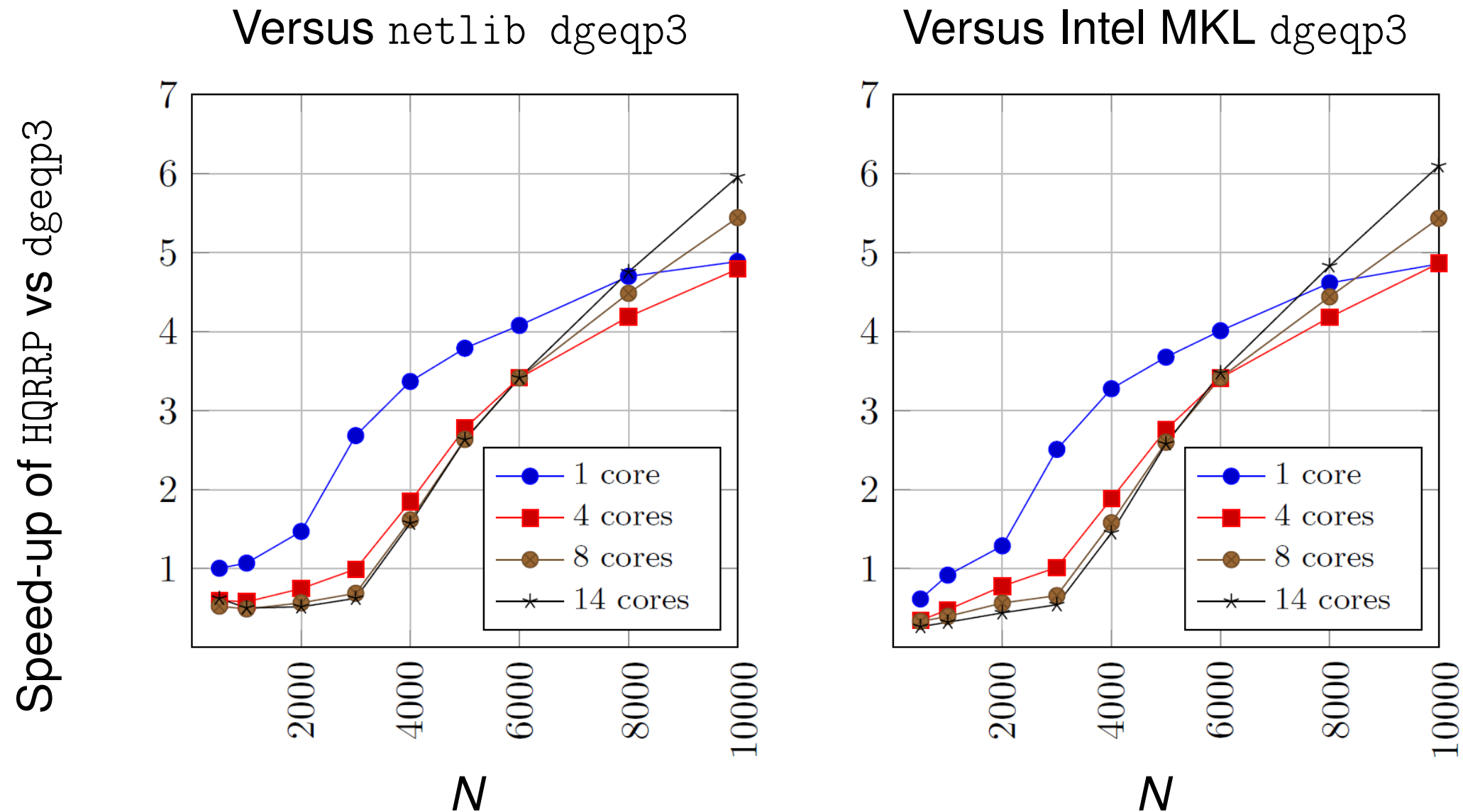
Everything is implemented in Matlab. The “full qr” line refers to Matlab built in qr.  
 CPU = Intel Xeon E-1660 (6 cores, 3.3GHz). GPU = Tesla K40c (2880 cores, 12GB).  
*Caveat: Matlab overhead makes column-pivoted QR slower than it could be.*

Time for compression of  $n \times n$  matrix.  $k=200$   $kstep=40$



Everything is implemented in Matlab. The “full qr” line refers to Matlab built in qr.  
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# Randomized algorithms for FULL factorizations: Column pivoted QR



Speedup attained by our randomized algorithm *HQRRP* for computing a full column pivoted QR factorization of an  $N \times N$  matrix. The speed-up is measured versus LAPACK's faster routine *dgeqp3* as implemented in Netlib (left) and Intel's MKL (right). Our implementation was done in C, and was executed on an Intel Xeon E5-2695. Joint work with G. Quintana-Ortí, N. Heavner, and R. van de Geijn. Available at: <https://github.com/flame/hqrrp/>